CS156: The Calculus of Computation

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Chapter 7: Quantified Linear Arithmetic

Quantifier Elimination (QE)

Algorithm for elimination of all quantifiers of formula F until quantifier-free formula (qff) G that is equivalent to F remains

Note: Could be enough if F is equisatisfiable to G, that is F is satisfiable iff G is satisfiable

A theory T admits quantifier elimination iff

there is an algorithm that given Σ-formula F returns

a quantifier-free Σ-formula G that is T-equivalent to F.

Example: $\exists x.\ 2x = y$

For $\Sigma_{\mathbb{Q}}$-formula

$$F:\ \exists x.\ 2x = y,$$

quantifier-free $T_{\mathbb{Q}}$-equivalent $\Sigma_{\mathbb{Q}}$-formula is

$$G:\ \top$$

For $\Sigma_{\mathbb{Z}}$-formula

$$F:\ \exists x.\ 2x = y,$$

there is no quantifier-free $T_{\mathbb{Z}}$-equivalent $\Sigma_{\mathbb{Z}}$-formula.

Let $\widehat{T_{\mathbb{Z}}}$ be $T_{\mathbb{Z}}$ with divisibility predicates $|$.
For $\widehat{\Sigma_{\mathbb{Z}}}$-formula

$$F:\ \exists x.\ 2x = y,$$

a quantifier-free $\widehat{T_{\mathbb{Z}}}$-equivalent $\widehat{\Sigma_{\mathbb{Z}}}$-formula is

$$G:\ 2 \mid y.$$

About QE Algorithm

In developing a QE algorithm for theory T, we need only consider formulae of the form

$$\exists x.\ F$$

for quantifier-free F.

<u>Example</u>: For Σ-formula

$$\begin{aligned}
G_1 &: \exists x.\ \forall y.\ \underbrace{\exists z.\ F_1[x, y, z]}_{F_2[x,y]} \\
G_2 &: \exists x.\ \forall y.\ F_2[x, y] \\
G_3 &: \exists x.\ \neg \underbrace{\exists y.\ \neg F_2[x, y]}_{F_3[x]} \\
G_4 &: \underbrace{\exists x.\ \neg F_3[x]}_{F_4} \\
G_5 &: F_4
\end{aligned}$$

G_5 is quantifier-free and T-equivalent to G_1

Quantifier Elimination for $T_{\mathbb{Z}}$

$\Sigma_{\mathbb{Z}} : \{\ldots, -2, -1, 0, 1, 2, \ldots, -3\cdot, -2\cdot, 2\cdot, 3\cdot, \ldots, +, -, =, <\}$

Lemma:

Given quantifier-free $\Sigma_{\mathbb{Z}}$-formula $F[y]$ s.t. $\text{free}(F[y]) = \{y\}$.
S represents the set of integers

$$S : \{n \in \mathbb{Z} \;:\; F[n] \text{ is } T_{\mathbb{Z}}\text{-valid}\} \;.$$

Either $S \cap \mathbb{Z}^+$ or $\mathbb{Z}^+ \setminus S$ is finite.
Note: $\mathbb{Z}^+$ is the set of positive integers.

Example: $\Sigma_{\mathbb{Z}}$-formula $\quad F[y] : \exists x.\ 2x = y$

S: even integers

$S \cap \mathbb{Z}^+$: positive even integers — infinite
$\mathbb{Z}^+ \setminus S$: positive odd integers — infinite

Therefore, by the lemma, there is no quantifier-free $T_{\mathbb{Z}}$-formula that is $T_{\mathbb{Z}}$-equivalent to $F[y]$.

Thus, $T_{\mathbb{Z}}$ does not admit QE.

Augmented theory $\widehat{T_{\mathbb{Z}}}$

$\widehat{\Sigma_{\mathbb{Z}}}$: $\Sigma_{\mathbb{Z}}$ with countable number of unary divisibility predicates

$$k \mid \cdot \quad \text{for } k \in \mathbb{Z}^+$$

Intended interpretations:

$k \mid x$ holds iff k divides x without any remainder

Example:

$$x > 1 \;\wedge\; y > 1 \;\wedge\; 2 \mid x + y$$

is satisfiable (choose $x = 2, y = 2$).

$$\neg(2 \mid x) \;\wedge\; 4 \mid x$$

is not satisfiable.

Axioms of $\widehat{T_{\mathbb{Z}}}$: axioms of $T_{\mathbb{Z}}$ with additional countable set of axioms

$$\forall x.\; k \mid x \;\leftrightarrow\; \exists y.\; x = ky \quad \text{for } k \in \mathbb{Z}^+$$

$\widehat{T_{\mathbb{Z}}}$ admits QE (Cooper's method)

Algorithm: Given $\widehat{\Sigma_{\mathbb{Z}}}$-formula

$$\exists x.\ F[x]\ ,$$

where F is quantifier-free, construct quantifier-free $\widehat{\Sigma_{\mathbb{Z}}}$-formula that is equivalent to $\exists x.\ F[x]$.

1. Put F[x] into Negation Normal Form (NNF).
2. Normalize literals: $s < t$, $k|t$, or $\neg(k|t)$.
3. Put x in $s < t$ on one side: $hx < t$ or $s < hx$.
4. Replace hx with x' without a factor.
5. Replace $F[x']$ by $\bigvee F[j]$ for finitely many j.

Cooper's Method: Step 1

Put $F[x]$ in Negation Normal Form (NNF) $F_1[x]$, so that $\exists x.\ F_1[x]$

- has negations only in literals (only $\wedge$, $\vee$)
- is $\widehat{T_{\mathbb{Z}}}$-equivalent to $\exists x.\ F[x]$

Example:

$$\exists x.\ \neg(x - 6 < z - x \ \wedge\ 4 \mid 5x + 1 \ \rightarrow\ 3x < y)$$

is equivalent to

$$\exists x.\ x - 6 < z - x \ \wedge\ 4 \mid 5x + 1 \ \wedge\ \neg(3x < y)$$

Note:

$$\neg(A \wedge B \rightarrow C) \Leftrightarrow (A \wedge B \wedge \neg C)$$

Cooper's Method: Step 2

Replace (left to right)

$$\begin{array}{rcl} s = t & \Leftrightarrow & s < t + 1 \ \wedge \ t < s + 1 \\ \neg(s = t) & \Leftrightarrow & s < t \ \vee \ t < s \\ \neg(s < t) & \Leftrightarrow & t < s + 1 \end{array}$$

The output $\exists x.\ F_2[x]$ contains only literals of form

$$s < t\ , \quad k \mid t\ , \quad \text{or} \quad \neg(k \mid t)\ ,$$

where s, t are $\widehat{T_{\mathbb{Z}}}$-terms and $k \in \mathbb{Z}^+$.

Example:

$$\neg(x < y) \ \wedge \ \neg(x = y + 3)$$
$$\Downarrow$$
$$y < x + 1 \ \wedge \ (x < y + 3 \ \vee \ y + 3 < x)$$

Cooper's Method: Step 3

Collect terms containing x so that literals have the form

$$hx < t \;, \quad t < hx \;, \quad k \mid hx + t \;, \quad \text{or} \quad \neg(k \mid hx + t) \;,$$

where t is a term (does not contain x) and $h, k \in \mathbb{Z}^+$. The output is the formula $\exists x.\ F_3[x]$, which is $\widehat{T_{\mathbb{Z}}}$-equivalent to $\exists x.\ F[x]$.

Example:

$$x + x + y < z + 3z + 2y - 4x \quad\Downarrow\quad 6x < 4z + y$$

$$5 | -7x + t \quad\Downarrow\quad 5 | 7x - t$$

Cooper's Method: Step 4 I

Let

$$\delta' = \text{lcm}\{h \ : \ h \text{ is a coefficient of } x \text{ in } F_3[x]\} \ ,$$

where lcm is the least common multiple. Multiply atoms in $F_3[x]$ by constants so that δ' is the coefficient of x everywhere:

$$\begin{array}{rcll} hx < t & \Leftrightarrow & \delta' x < h' t & \text{where} \quad h' h = \delta' \\ t < hx & \Leftrightarrow & h' t < \delta' x & \text{where} \quad h' h = \delta' \\ k \mid hx + t & \Leftrightarrow & h' k \mid \delta' x + h' t & \text{where} \quad h' h = \delta' \\ \neg(k \mid hx + t) & \Leftrightarrow & \neg(h' k \mid \delta' x + h' t) & \text{where} \quad h' h = \delta' \end{array}$$

The result $\exists x.\ F_3'[x]$, in which all occurrences of x in $F_3'[x]$ are in terms $\delta' x$.

Replace $\delta' x$ terms in F_3' with a fresh variable x' to form

$$F_3'' \ : \ F_3\{\delta' x \mapsto x'\}$$

Cooper's Method: Step 4 II

Finally, construct

$$\exists x'. \ \underbrace{F_3''[x'] \ \wedge \ \delta' \mid x'}_{F_4[x']}$$

$\exists x'.F_4[x']$ is equivalent to $\exists x.\ F[x]$ and each literal of $F_4[x']$ has one of the forms:

(A) $x' < t$

(B) $t < x'$

(C) $k \mid x' + t$

(D) $\neg(k \mid x' + t)$

where t is a term that does not contain x', and $k \in \mathbb{Z}^+$.

Cooper's Method: Step 4 III

Example: $\widehat{T_{\mathbb{Z}}}$-formula

$$\exists x.\ \underbrace{3x + 1 > y \ \wedge\ 2x - 6 < z \ \wedge\ 4 \mid 5x + 1}_{F[x]}$$

After step 3:

$$\exists x.\ \underbrace{2x < z + 6 \ \wedge\ y - 1 < 3x \ \wedge\ 4 \mid 5x + 1}_{F_3[x]}$$

Collecting coefficients of x (step 4):

$$\delta' = \text{lcm}(2, 3, 5) = 30$$

Multiply when necessary:

$$\exists x.\ 30x < 15z + 90 \ \wedge\ 10y - 10 < 30x \ \wedge\ 24 \mid 30x + 6$$

Cooper's Method: Step 4 IV

Multiply when necessary:

$$\exists x.\ 30x < 15z + 90 \ \land\ 10y - 10 < 30x \ \land\ 24 \mid 30x + 6$$

Replacing $30x$ with fresh x' and adding divisibility conjunct:

$$\exists x'.\ \underbrace{x' < 15z + 90 \ \land\ 10y - 10 < x' \ \land\ 24 \mid x' + 6 \ \land\ 30 \mid x'}_{F_4[x']}$$

$\exists x'.\ F_4[x']$ is equivalent to $\exists x.\ F[x]$.

Cooper's Method: Step 5

Construct left infinite projection $F_{-\infty}[x']$ of $F_4[x']$ by

(A) replacing literals $x' < t$ by $\top$

(B) replacing literals $t < x'$ by $\bot$

Idea: very small numbers satisfy (A) literals but not (B) literals

Let

$$\delta = \text{lcm} \left\{ \begin{array}{l} k \text{ of (C) literals } k \mid x' + t \\ k \text{ of (D) literals } \neg(k \mid x' + t) \end{array} \right\}$$

and B be the set of terms t appearing in (B) literals of $F_4[x']$.

Construct

$$F_5 : \bigvee_{j=1}^{\delta} F_{-\infty}[j] \lor \bigvee_{j=1}^{\delta} \bigvee_{t \in B} F_4[t + j] .$$

F_5 is quantifier-free and $\widehat{T_{\mathbb{Z}}}$-equivalent to $\exists x.\ F[x]$.

Intuition of Step 5 I

Property (Periodicity)

if $m \mid \delta$
then $m \mid n$ iff $m \mid n + \lambda\delta$ for all $\lambda \in \mathbb{Z}$

That is, $m \mid \cdot$ cannot distinguish between $m \mid n$ and $m \mid n + \lambda\delta$.

By the choice of δ (lcm of the k's) — no $\mid$ literal in F_5 can distinguish between n and $n + \lambda\delta$, for any $\lambda \in \mathbb{Z}$.

$$F_5 : \bigvee_{j=1}^{\delta} F_{-\infty}[j] \ \vee \ \bigvee_{j=1}^{\delta} \bigvee_{t \in B} F_4[t + j]$$

Intuition of Step 5 II

- left disjunct $\bigvee_{j=1}^{\delta} F_{-\infty}[j]$:

  Contains only $\mid$ literals
  Asserts: no least $n \in \mathbb{Z}$ s.t. $F_4[n]$.

  For if there exists n satisfying $F_{-\infty}$,
  then every $n - \lambda\delta$, for $\lambda \in \mathbb{Z}^+$, also satisfies $F_{-\infty}$

- right disjunct $\bigvee_{j=1}^{\delta} \bigvee_{t \in B} F_4[t + j]$:

  Asserts: There is least $n \in \mathbb{Z}$ s.t. $F_4[n]$.

  For let $t^* = \{\text{largest } t \mid t < x' \text{ in (B)}\}$.
  If $n \in \mathbb{Z}$ is s.t. $F_4[n]$, then

  $$\exists j (1 \leq j \leq \delta).\ t^* + j \leq n \ \wedge\ F_4[t^* + j]$$

  In other words,
  if there is a solution,
  then one must appear in δ interval to the right of t^*

Example of Step 5 I

$$\exists x.\ \underbrace{3x + 1 > y \ \land\ 2x - 6 < z \ \land\ 4 \mid 5x + 1}_{F[x]}$$

$$\Downarrow$$

$$\exists x'.\ \underbrace{x' < 15z + 90 \ \land\ 10y - 10 < x' \ \land\ 24 \mid x' + 6 \ \land\ 30 \mid x'}_{F_4[x']}$$

By step 5,

$$F_{-\infty}[x'] :\ \top \ \land\ \bot \ \land\ 24 \mid x' + 6 \ \land\ 30 \mid x' \ ,$$

which simplifies to $\bot$.

Example of Step 5 II

Compute

$$\delta = \text{lcm}\{24, 30\} = 120 \quad \text{and} \quad B = \{10y - 10\} \ .$$

Then replacing x' by $10y - 10 + j$ in $F_4[x']$ produces

$$F_5 : \bigvee_{j=1}^{120} \left[\begin{array}{l} 10y - 10 + j < 15z + 90 \ \wedge \ 10y - 10 < 10y - 10 + j \\ \ \wedge \ 24 \mid 10y - 10 + j + 6 \ \wedge \ 30 \mid 10y - 10 + j \end{array} \right]$$

which simplifies to

$$F_5 : \bigvee_{j=1}^{120} \left[\begin{array}{l} 10y + j < 15z + 100 \ \wedge \ \cancel{0 < j} \\ \ \wedge \ 24 \mid 10y + j - 4 \ \wedge \ 30 \mid 10y - 10 + j \end{array} \right] .$$

F_5 is quantifier-free and $\widehat{T_{\mathbb{Z}}}$-equivalent to $\exists x.\ F[x]$.

Cooper's Method: Example I

$$\underbrace{\exists x.\ (3x + 1 < 10\ \lor\ 7x - 6 > 7)\ \land\ 2 \mid x}_{F[x]}$$

Isolate x terms

$$\exists x.\ (3x < 9\ \lor\ 13 < 7x)\ \land\ 2 \mid x\ ,$$

so

$$\delta' = \text{lcm}\{3, 7, 1\} = 21\ .$$

After multiplying coefficients by proper constants,

$$\exists x.\ (21x < 63\ \lor\ 39 < 21x)\ \land\ 42 \mid 21x\ ,$$

we replace $21x$ by x':

$$\exists x'.\ \underbrace{(x' < 63\ \lor\ 39 < x')\ \land\ 42 \mid x'\ \land\ 21 \mid x'}_{F_4[x']}\ .$$

Cooper's Method: Example II

Then

$$F_{-\infty}[x'] : \ (\top \ \vee \ \bot) \ \wedge \ 42 \mid x' \ \wedge \ 21 \mid x' \ ,$$

or, simplifying,

$$F_{-\infty}[x'] : \ 42 \mid x' \ \wedge \ 21 \mid x' \ .$$

Finally,

$$\delta = \text{lcm}\{21, 42\} = 42 \quad \text{and} \quad B = \{39\} \ ,$$

so F_5 :

$$\bigvee_{j=1}^{42} (42 \mid j \ \wedge \ 21 \mid j) \quad \vee$$

$$\bigvee_{j=1}^{42} ((39 + j < 63 \ \vee \ 39 < 39 + j) \ \wedge \ 42 \mid 39 + j \ \wedge \ 21 \mid 39 + j) \ .$$

Since $42 \mid 42$ and $21 \mid 42$, the left main disjunct simplifies to $\top$, so that F_5 is $\widehat{T_{\mathbb{Z}}}$-equivalent to $\top$. Thus, $\exists x.\ F[x]$ is $\widehat{T_{\mathbb{Z}}}$-valid.

Cooper's Method: Example I

$$\exists x.\ \underbrace{2x = y}_{F[x]}$$

Rewriting

$$\exists x.\ \underbrace{2x < y + 1 \ \wedge \ y - 1 < 2x}_{F_3[x]}$$

Then

$$\delta' = \text{lcm}\{2, 2\} = 2\ ,$$

so by Step 4

$$\exists x'.\ \underbrace{x' < y + 1 \ \wedge \ y - 1 < x' \ \wedge \ 2 \mid x'}_{F_4[x']}$$

$F_{-\infty}$ produces $\perp$.

Cooper's Method: Example II

However,

$$\delta = \text{lcm}\{2\} = 2 \quad \text{and} \quad B = \{y - 1\} \ ,$$

so

$$F_5 : \ \bigvee_{j=1}^{2} (y - 1 + j < y + 1 \ \wedge \ y - 1 < y - 1 + j \ \wedge \ 2 \mid y - 1 + j)$$

Simplifying,

$$F_5 : \ \bigvee_{j=1}^{2} (j < 2 \ \wedge \ 0 < j \ \wedge \ 2 \mid y - 1 + j)$$

and then

$$F_5 : \ 2 \mid y \ ,$$

which is quantifier-free and $\widehat{T_{\mathbb{Z}}}$-equivalent to $\exists x.\ F[x]$.

Improvement: Symmetric Elimination

In step 5, if there are fewer

(A) literals $x' < t$

than

(B) literals $t < x'$,

construct the right infinite projection $F_{+\infty}[x']$ from $F_4[x']$ by replacing

(A) literal $x' < t$ by $\perp$

than

(B) literal $t < x'$ by $\top$

Then right elimination.

$$F_5 : \bigvee_{j=1}^{\delta} F_{+\infty}[-j] \ \vee \ \bigvee_{j=1}^{\delta} \bigvee_{t \in A} F_4[t - j] \ .$$

Improvement: Eliminating Blocks of Quantifiers I

Given

$$\exists x_1. \cdots \exists x_n. \ F[x_1, \ldots, x_n]$$

where F quantifier-free.
Eliminating x_n (left elimination) produces

$$G_1 : \quad \exists x_1. \cdots \exists x_{n-1}. \quad \bigvee_{j=1}^{\delta} F_{-\infty}[x_1, \ldots, x_{n-1}, j] \ \vee \\ \bigvee_{j=1}^{\delta} \bigvee_{t \in B} F_4[x_1, \ldots, x_{n-1}, t + j]$$

which is equivalent to

$$G_2 : \quad \bigvee_{j=1}^{\delta} \exists x_1. \cdots \exists x_{n-1}. \ F_{-\infty}[x_1, \ldots, x_{n-1}, j] \ \vee \\ \bigvee_{j=1}^{\delta} \bigvee_{t \in B} \exists x_1. \cdots \exists x_{n-1}. \ F_4[x_1, \ldots, x_{n-1}, t + j]$$

Improvement: Eliminating Blocks of Quantifiers II

Treat j as a free variable and examine only $1 + |B|$ formulae

- $\exists x_1. \cdots \exists x_{n-1}. F_{-\infty}[x_1, \ldots, x_{n-1}, j]$
- $\exists x_1. \cdots \exists x_{n-1}. F_4[x_1, \ldots, x_{n-1}, t + j]$ for each $t \in B$

Example I

$$F: \ \exists y.\ \exists x.\ x < -2 \ \wedge\ 1 - 5y < x \ \wedge\ 1 + y < 13x$$

Since $\delta' = \text{lcm}\{1, 13\} = 13$

$$\exists y.\ \exists x.\ 13x < -26 \ \wedge\ 13 - 65y < 13x \ \wedge\ 1 + y < 13x$$

Then

$$\exists y.\ \exists x'.\ x' < -26 \ \wedge\ 13 - 65y < x' \ \wedge\ 1 + y < x' \ \wedge\ 13 \mid x'$$

There is one (A) literal $x' < \ldots$ and two (B) literals $\ldots < x'$, we use right elimination.

$$F_{+\infty} = \bot \qquad \delta = \{13\} = 13 \qquad A = \{-26\}$$

$$F': \ \exists y.\ \bigvee_{j=1}^{13} \left[\begin{array}{c} -26 - j < -26 \ \wedge\ 13 - 65y < -26 - j \\ \wedge\ 1 + y < -26 - j \ \wedge\ 13 \mid \ -26 - j \end{array} \right]$$

Example II

Commute

$$G[j] : \bigvee_{j=1}^{13} \underbrace{\exists y.\, j > 0 \ \wedge\ 39 + j < 65y \ \wedge\ y < -27 - j \ \wedge\ 13 \mid -26 - j}_{H[j]}$$

Treating j as free variable (and removing $j > 0$), apply QE to

$$H[j] :\ \exists y.\, 39 + j < 65y \ \wedge\ y < -27 - j \ \wedge\ 13 \mid -26 - j$$

Simplify...

$$H'[j] :\ \bigvee_{k=1}^{65} (k < -1794 - 66j \ \wedge\ 13 \mid -26 - j \ \wedge\ 65 \mid 39 + j + k)$$

Replace $H[j]$ with $H'[j]$ in $G[j]$

Example III

$$F'' : \bigvee_{j=1}^{13} \bigvee_{k=1}^{65} (k < -1794 - 66j \wedge 13 \mid -26 - j \wedge 65 \mid 39 + j + k)$$

$$\uparrow \qquad\qquad \uparrow$$

$$j = 13 \qquad\qquad k = 13$$

simplified to

$$13 < -1794 - 66 \cdot 13$$

$$\bot$$

This qff formula is $\widehat{T_{\mathbb{Z}}}$-equivalent to F.

Quantifier Elimination over Rationals

$$\Sigma_{\mathbb{Q}} : \{0, 1, +, -, =, \geq\}$$

Recall: we use $>$ instead of $\geq$, as

$$x \geq y \Leftrightarrow x > y \lor x = y \qquad x > y \Leftrightarrow x \geq y \land \neg(x = y) .$$

Ferrante & Rackoff's Method

Given a $\Sigma_{\mathbb{Q}}$-formula $\exists x.\ F[x]$, where $F[x]$ is quantifier-free, generate quantifier-free formula F_4 (four steps) s.t.

$$F_4 \text{ is } \Sigma_{\mathbb{Q}}\text{-equivalent to } \exists x.\ F[x]$$

by

1. putting $F[x]$ in NNF,
2. replacing negated literals,
3. solving literals such that x appears isolated on one side, and
4. taking finite disjunction $\bigvee_t F[t]$.

Ferrante & Rackoff's Method: Steps 1 and 2

Step 1: Put $F[x]$ in NNF. The result is $\exists x.\ F_1[x]$.

Step 2: Replace literals (left to right)

$$\neg(s < t) \;\Leftrightarrow\; t < s \;\vee\; t = s$$
$$\neg(s = t) \;\Leftrightarrow\; t < s \;\vee\; t > s$$

The result $\exists x.\ F_2[x]$ does not contain negations.

Ferrante & Rackoff's Method: Step 3

Solve for x in each atom of $F_2[x]$, e.g.,

$$t_1 < cx + t_2 \qquad \Rightarrow \qquad \frac{t_1 - t_2}{c} < x$$

where $c \in \mathbb{Z} - \{0\}$.

All atoms in the result $\exists x.\ F_3[x]$ have form

(A) $x < t$

(B) $t < x$

(C) $x = t$

where t is a term that does not contain x.

Ferrante & Rackoff's Method: Step 4 I

Construct from $F_3[x]$

- left infinite projection $F_{-\infty}$ by replacing
  (A) atoms $x < t$ by $\top$
  (B) atoms $t < x$ by $\bot$
  (C) atoms $x = t$ by $\bot$
- right infinite projection $F_{+\infty}$ by replacing
  (A) atoms $x < t$ by $\bot$
  (B) atoms $t < x$ by $\top$
  (C) atoms $x = t$ by $\bot$

Let S be the set of t terms from (A), (B), (C) atoms.
Construct the final

$$F_4 : \; F_{-\infty} \;\vee\; F_{+\infty} \;\vee\; \bigvee_{s,t \in S} F_3\left[\frac{s+t}{2}\right] ,$$

which is $T_{\mathbb{Q}}$-equivalent to $\exists x.\ F[x]$.

Ferrante & Rackoff's Method: Step 4 II

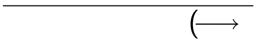
- $F_{-\infty}$ captures the case when small $x \in \mathbb{Q}$ satisfy $F_3[x]$
- $F_{+\infty}$ captures the case when large $x \in \mathbb{Q}$ satisfy $F_3[x]$
- last disjunct: for $s, t \in S$
  if $s \equiv t$, check whether $s \in S$ satisfies $F_3[s]$
  if $s \not\equiv t$, in any $T_{\mathbb{Q}}$-interpretation,
  - $|S| - 1$ pairs $s, t \in S$ are adjacent. For each such pair, (s, t) is an interval in which no other $s' \in S$ lies.
  - Since $\frac{s+t}{2}$ represents the whole interval (s, t), simply check $F_3[\frac{s+t}{2}]$.

Ferrante & Rackoff's Method: Intuition

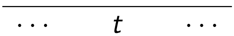
Step 4 says that four cases are possible:

1. There is a left open interval s.t. all elements satisfy $F(x)$.

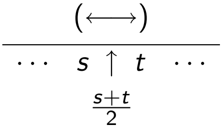
2. There is a right open interval s.t. all elements satisfy $F(x)$.



3. Some term t satisfies $F(x)$.



4. There is an open interval between two s, t terms such that every element satisfies $F(x)$.



Correctness of Step 4 I

Theorem

Let

$$F_4 : \; F_{-\infty} \;\vee\; F_{+\infty} \;\vee\; \bigvee_{s,t \in S} F_3\left[\frac{s+t}{2}\right] \;,$$

be the formula constructed from $\exists x.\; F_3[x]$ *as in Step 4. Then* $\exists x.\; F_3[x] \;\Leftrightarrow\; F_4$.

<u>Proof</u>:

$\Leftarrow$ If F_4 is true, then $F_{-\infty}$, F_∞ or $F_3[\frac{s+t}{2}]$ is true.
If $F_3[\frac{s+t}{2}]$ is true, then obviously $\exists x.\; F_3[x]$ is true.
If $F_{-\infty}$ is true, choose some small $x, x < t$ for all $t \in S$.
Then $F_3[x]$ is true.
If $F_{+\infty}$ is true, choose some big $x, x > t$ for all $t \in S$.
Then $F_3[x]$ is true.

Correctness of Step 4 II

$\Rightarrow$ If $I \models \exists x.\ F_3[x]$ then there is value v such that

$$I \models F_3[v].$$

If v $< \alpha_I[t]$ for all $t \in S$, then $I \models F_{-\infty}$.
If v $> \alpha_I[t]$ for all $t \in S$, then $I \models F_{+\infty}$.
If v $= \alpha_I[t]$ for some $t \in S$, then $I \models F[\frac{t+t}{2}]$.

Otherwise choose largest $s \in S$ with $\alpha_I[s] <$ v and smallest $t \in S$ with $\alpha_I[t] >$ v.

Since no atom of F_3 can distinguish between values in interval (s, t),

$$I \models F_3[\mathrm{v}] \quad \text{iff} \quad I \models F_3\left[\frac{s+t}{2}\right].$$

Hence, $I \models F[\frac{s+t}{2}]$. In all cases $I \models F_4$.

Ferrante & Rackoff's Method: Example I

$\Sigma_{\mathbb{Q}}$-formula

$$\exists x.\ \underbrace{3x + 1 < 10\ \wedge\ 7x - 6 > 7}_{F[x]}$$

Solving for x

$$\exists x.\ \underbrace{x < 3\ \wedge\ x > \frac{13}{7}}_{F_3[x]}$$

Step 4: $x > \frac{13}{7}$ in (B) $\Rightarrow$ $F_{-\infty} = \bot$
$x < 3$ in (A) $\Rightarrow$ $F_{+\infty} = \bot$

$$F_4 :\ \bigvee_{s,t \in S} \underbrace{\left(\frac{s+t}{2} < 3\ \wedge\ \frac{s+t}{2} > \frac{13}{7}\right)}_{F_3[\frac{s+t}{2}]}$$

Ferrante & Rackoff's Method: Example II

$S = \{3, \frac{13}{7}\} \quad \Rightarrow$

$$F_3\left[\frac{3+3}{2}\right] = \bot \qquad F_3\left[\frac{\frac{13}{7}+\frac{13}{7}}{2}\right] = \bot$$

$$F_3\left[\frac{\frac{13}{7}+3}{2}\right] : \ \frac{\frac{13}{7}+3}{2} < 3 \ \wedge \ \frac{\frac{13}{7}+3}{2} > \frac{13}{7} = \top$$

$$F_4 : \quad \bot \vee \cdots \vee \bot \vee \top = \top$$

Thus, $F_4 : \top$ is $T_{\mathbb{Q}}$-equivalent to $\exists x.\ F[x]$,

so $\exists x.\ F[x]$ is $T_{\mathbb{Q}}$-valid.

Example

$$\exists x.\ \underbrace{2x > y \ \wedge\ 3x < z}_{F[x]}$$

Solving for x

$$\exists x.\ \underbrace{x > \frac{y}{2} \ \wedge\ x < \frac{z}{3}}_{F_3[x]}$$

Step 4: $F_{-\infty} = \bot$, $F_{+\infty} = \bot$, $F_3[\frac{y}{2}] = \bot$ and $F_3[\frac{z}{3}] = \bot$.

$$F_4 :\ \frac{\frac{y}{2} + \frac{z}{3}}{2} > \frac{y}{2} \ \wedge\ \frac{\frac{y}{2} + \frac{z}{3}}{2} < \frac{z}{3}$$

which simplifies to:

$$F_4 :\ 2z > 3y$$

F_4 is $T_{\mathbb{Q}}$-equivalent to $\exists x.\ F[x]$.