CS256/Spring 2008 — Lecture #7

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Strengthening vs. Incremental Proof

Comparing the Strategies

We want to prove $\Box\, q$, but q is not inductive.

We have two options:

1. Strengthening

   Strengthen it to $q \wedge \varphi$.

   Prove $\Box(q \wedge \varphi)$ and deduce $\Box\, q$.

2. Incremental

   First prove $\Box\, \varphi$ and then prove $\Box\, q$ relative to φ.

Resulting verification conditions:

1

I1. $\Theta \;\rightarrow\; q \wedge \varphi$

I2. $\{q \wedge \varphi\}\ \mathcal{T}\ \{q \wedge \varphi\}$

2

$$\frac{\text{I1'.}\ \ \Theta \rightarrow \varphi \qquad \text{I2'.}\ \ \{\varphi\}\ \mathcal{T}\ \{\varphi\}}{\Box\, \varphi} \qquad \frac{\text{I1''.}\ \ \Theta \rightarrow q \qquad \text{I2''.}\ \ \{q \wedge \varphi\}\ \mathcal{T}\ \{q\}}{\Box\, q}$$

Strengthening vs. Incremental Proof (Con't)

- $\boxed{1}$ is strictly more powerful than $\boxed{2}$.

  $\boxed{2}$ implies $\boxed{1}$ since

$$\left[\begin{array}{c}\underbrace{\rho_\tau \land \varphi \ \rightarrow \varphi'}_{\text{I2'}} \\ \\ \underbrace{\rho_\tau \land q \land \varphi \rightarrow q'}_{\text{I2''}}\end{array}\right] \rightarrow \ [\underbrace{\rho_\tau \land q \land \varphi \ \rightarrow \ q' \land \varphi'}_{\text{I2}}]$$

- In practice, $\boxed{2}$ is often more useful than $\boxed{1}$
  - allows breaking down the proof in more manageable pieces
  - smaller verification conditions
  - more intuitive

Strengthening vs. Incremental Proof (Con't)

`Example:`

$$\begin{array}{l}\textbf{local } x\text{: } \textbf{integer where } x = 1 \\ \ell_0\text{: } \textbf{loop forever do} \\ \quad \left[\ \ell_1 : \ x := x + 1 \ \right]\end{array}$$

Show $q_1\colon \quad at\_\ell_0 \ \rightarrow \ x > 0$

$q_2\colon \quad at\_\ell_1 \ \rightarrow \ x > 0$

- both are P-valid
- neither of them is inductive
- but $q_1 \wedge q_2$ is inductive!

Combining the Strategies

Rule INC-INV**:** (incremental invariance)

For assertions q,φ,χ_1,. . . ,χ_k

$$\text{I0.} \quad P \vDash \Box\, \chi_1, \ldots, \Box\, \chi_k$$

$$\text{I1.} \quad P \Vdash (\bigwedge_{i=1}^{k} \chi_i) \wedge \varphi \rightarrow q$$

$$\text{I2.} \quad P \Vdash \Theta \rightarrow \varphi$$

$$\text{I3.} \quad P \Vdash \Big\{(\bigwedge_{i=1}^{k} \chi_i) \wedge \varphi\Big\}\ \mathcal{T}\ \{\varphi\}$$

$$P \vDash \Box\, q$$

If φ satisfies I2 and I3, we say that

"φ is inductive relative to $\chi_1, \ldots \chi_k$"

Combining the Strategies (Con't)

Note that Θ must be stronger than all the χ_i's (i.e., $P \Vdash \Theta \rightarrow \chi_i$) and so

$$P \Vdash \left(\bigwedge_{i=1}^{k} \chi_i \right) \wedge \Theta \rightarrow \varphi \quad \text{iff} \quad P \Vdash \Theta \rightarrow \varphi$$

From now on, we usually omit "$P \models$" and "$P \Vdash$".

Detecting Trivial Verification Conditions

$\{\varphi\}\ \mathcal{T}\ \{\varphi\}$ – Don't check every $\tau \in \mathcal{T}$.

- Ignore $\{\varphi\}\ \tau_I\ \{\varphi\}$ – always true
- Ignore $\{\varphi\}\ \tau\ \{\varphi\}$
  if τ does not modify any variable in φ
- For $\{\varphi\}\ \tau\ \{\varphi\}$ where φ: $p \rightarrow q$

$$\rho_\tau \ \wedge\ \underbrace{p \rightarrow q}_{\varphi}\ \rightarrow\ \underbrace{p' \rightarrow q'}_{\varphi'}$$

  Consider only τ's that
  validate p or falsify q

Finding Inductive Assertions

Two methods:

1. Bottom-up:
    - based on the program text only
    - algorithmic
    - guaranteed to produce an inductive invariant

2. Top-down:
    - guided by the property we want to prove
    - heuristic
    - not guaranteed to produce an inductive invariant

Finding Inductive Assertions
Bottom-Up Approach

- Transition-validated assertions:

ℓ_1: [**while** c **do** S]; ℓ_2: $\boxed{at\_\ell_2 \rightarrow \neg c}$

if no statement parallel to ℓ_2 can modify variables in c

ℓ_1: $y := e$; ℓ_2: $\boxed{at\_\ell_2 \rightarrow y = e}$

if no statement parallel to ℓ_2 can modify y or variables occurring in e and if y does not occur in e.

Bottom-Up Approach (Con't)

- single variable assertions

$$y = 1$$

$$\begin{bmatrix} \textbf{loop forever do} \\ \quad \begin{bmatrix} \cdots \\ \textbf{request } y \\ \cdots \\ \textbf{release } y \end{bmatrix} \end{bmatrix}$$

$$\boxed{y \geq 0}$$

$$s = 1$$

$$\begin{bmatrix} \cdots \\ s := 1 \\ \cdots \end{bmatrix} \; || \; \begin{bmatrix} \cdots \\ s := 2 \\ \cdots \end{bmatrix}$$

$$\boxed{s = 1 \;\vee\; s = 2}$$

where no other statement modifies s

`Example:` Program SQUARE-ROOT

Fig. 1.11

$$\boxed{at\_\ell_2 \;\rightarrow\; z^2 \leq x < (z+1)^2}$$

<u>Intuitive argument:</u>

$$z \;=\; 0, 1, \ldots, n$$

$$u \;=\; 1, 3, \ldots, 2n+1$$

$$w \;=\; \underbrace{1 + 3 + \ldots + (2n{+}1)}_{(n+1)^2} \;=\; (z+1)^2$$

first time $w > x$

$$x < (z+1)^2$$

last time $w \leq x$

$$z^2 \leq x$$

Thus at ℓ_2:

$$z^2 \leq x < (z+1)^2$$

Program SQUARE-ROOT

$$
\begin{array}{llll}
\textbf{in} & x\text{:} & \textbf{integer} & \textbf{where } x \geq 0 \\
\textbf{local} & u, w\text{:} & \textbf{integer} & \textbf{where } u = 1, w = 1 \\
\textbf{out} & z\text{:} & \textbf{integer} & \textbf{where } z = 0
\end{array}
$$

ℓ_0 : **while** $w \leq x$ **do**

  ℓ_1 : $(z,\ u,\ w)\ :=\ (z+1,\ u+2,\ w+u+2)$

ℓ_2 :

$$\rho_{\ell_0}\text{:}\quad \underbrace{move(\ell_0, \ell_1)\ \wedge\ w \leq x}_{\rho^{\mathrm{T}}_{\ell_0}}\quad \vee$$

$$\underbrace{move(\ell_0, \ell_2)\ \wedge\ w > x}_{\rho^{\mathrm{F}}_{\ell_0}}$$

$$
\begin{array}{lll}
\rho_{\ell_1}\text{:} & move(\ell_1, \ell_0) & \wedge \\
 & z' = z + 1 & \wedge \\
 & u' = u + 2 & \wedge \\
 & w' = w + u + 2 &
\end{array}
$$

Find $\psi_2\colon\ at\_\ell_2 \;\rightarrow\; x < (z+1)^2$

$$\begin{cases} z_0 = 0 \\ z_n = z_{n-1} + 1 \quad \text{for } n > 0 \end{cases}$$

$$\begin{cases} u_0 = 1 \\ u_n = u_{n-1} + 2 \quad \text{for } n > 0 \end{cases}$$

$$\begin{cases} w_0 = 1 \\ w_n = w_{n-1} + u_{n-1} + 2 \quad \text{for } n > 0 \end{cases}$$

- Step 1

$$\left.\begin{array}{ll} z_n = n & \text{for } n \geq 0 \\ u_n = 2n+1 & \text{for } n \geq 0 \end{array}\right\} \Rightarrow \begin{array}{l} u_n = 2z_n + 1 \\ \text{for } n \geq 0 \end{array}$$

$$\boxed{\varphi_1\colon\ u = 2z + 1}$$

- Step 2

$$\begin{cases} w_0 = 1 \\ w_n = w_{n-1} + (\overbrace{2(n-1)+1}^{u_{n-1}}) + 2 \\ \quad = w_{n-1} + (2n+1) \qquad\qquad \text{for } n \geq 0 \end{cases}$$

$$w_n = \sum_{k=0}^{n} (2k+1) = (n+1)^2 \quad \text{for } n \geq 0$$

$$w_n = (z_n + 1)^2 \quad \text{for } n \geq 0$$

$$\boxed{\varphi_2\text{: } w = (z+1)^2}$$

- Step3

$$\boxed{at\_\ell_2 \;\rightarrow\; x < w}$$

Therefore

$$\boxed{\psi_2\text{: } at\_\ell_2 \;\rightarrow\; x < (z+1)^2}$$

Construction of Linear Invariants

a limited class of invariants that can be constructed algorithmically

Definition: integer variable y is linear in P if

$$y' = y + c \qquad \text{for every } \rho_\tau$$

for some integer constant c.

`Example:` semaphore variables are linear

$$\underbrace{y' = y + 1}_{\textbf{release}} \qquad \underbrace{y' = y - 1}_{\textbf{request}} \qquad \underbrace{y' = y}_{\text{otherwise}}$$

Definition:

A linear invariant is of the form

$$\underbrace{\sum_{i=1}^{r} a_i \cdot y_i}_{\text{body}} + \underbrace{\sum_{\ell \in \mathcal{L}} b_\ell \cdot at\_\ell}_{\text{compensation expression}} = \underbrace{K}_{\text{constant}}$$

where

$a_i, \; b_\ell, \; K$ – integer constants.

$\mathcal{L}$ – set of all locations in P

$y_1, ..., y_r$ – all linear variables in P

`Example:` Program DOUBLE

$$\textbf{local } y\text{: } \textbf{integer where } y = 0$$

$$\begin{bmatrix} \ell_0\text{: } y := y+1 \\ \ell_1\text{:} \end{bmatrix} \quad \| \quad \begin{bmatrix} m_0\text{: } y := y+1 \\ m_1\text{:} \end{bmatrix}$$

linear variable: y

linear invariant:

$$\boxed{y + at\_\ell_0 + at\_m_0 \;=\; 2}$$

How are linear invariants constructed?

Our procedure guarantees that the generated assertions are P-invariants!

Assumption

Program $\ell_0^1\colon\ S_1 \parallel \ldots \parallel \ell_0^i\colon\ S_i \parallel \ \ldots\ \parallel \ell_0^m\colon\ S_m$

- no nested parallel statements. Therefore, all move expressions in all ρ_τ are of the form $move(\ell_i, \ell_j)$

- all linear variables y_i have a single initial value y_i^0

- every transition τ enabled on some P-accessible state

Increments

- $\Delta(y, \tau) = c$ if $\rho_\tau \rightarrow y' = y + c$

  therefore $\rho_\tau \rightarrow y' = y + \Delta(y, \tau)$

- $$\Delta(at\_\ell, \tau) \ = \begin{cases} 1 & \text{if } \ell = \ell_j \\ -1 & \text{if } \ell = \ell_i \\ 0 & \text{otherwise} \end{cases}$$

  if $\rho_\tau \rightarrow move(\ell_i, \ell_j)$

  therefore $\rho_\tau \rightarrow at'\_\ell = at\_\ell + \Delta(at\_\ell, \tau)$

Equations

Construct

$$\varphi: \quad \sum_{i=1}^{r} a_i \cdot y_i \; + \; \sum_{\ell \in \mathcal{L}} b_\ell \cdot at\_\ell \; = \; K$$

We obtain the values of the coefficients from a set of equations as follows:

(I) The invariant has to hold at the first state of every computation

Θ implies $y_i = y_i^0 \ (i = 1 \ldots r)$
and $\pi = \{\ell_0^1, \ldots, \ell_0^m\}$

and so we get

$$\boxed{\sum_{i=1}^{r} a_i \cdot y_i^0 \; + \; (b_{\ell_0^1} + \cdots + b_{\ell_0^m}) \; = \; K}$$

Equations (Con'd)

(T) the assertion has to be preserved by all transitions (we want it to be inductive):

$$\underbrace{\left(\sum_{i=1}^{r} a_i \cdot y_i \;+\; \sum_{\ell \in \mathcal{L}} b_\ell \cdot at\_\ell = K\right)}_{\varphi} \;\wedge\; \rho_\tau$$

$$\rightarrow\; \underbrace{\left(\sum_{i=1}^{r} a_i \cdot y_i' \;+\; \sum_{\ell \in \mathcal{L}} b_\ell \cdot at'\_\ell = K\right)}_{\varphi'}$$

or

$$\rho_\tau \;\rightarrow\; \sum_{i=1}^{r} a_i \cdot (y_i' - y_i) \;+\; \sum_{\ell \in \mathcal{L}} b_\ell \cdot (at'\_\ell - at\_\ell) \;=\; 0$$

resulting in the set of equations

$$\boxed{\sum_{i=1}^{r} a_i \cdot \Delta(y_i, \tau) \;+\; \sum_{\ell \in \mathcal{L}} b_\ell \cdot \Delta(at\_\ell, \tau) \;=\; 0}$$

for every transition $\tau \in \mathcal{T}$

`Example:` Program DOUBLE

$$\textbf{local } y\text{: } \textbf{integer where } y = 0$$

$$\begin{bmatrix} \ell_0\text{: } y := y + 1 \\ \ell_1\text{:} \end{bmatrix} \quad \| \quad \begin{bmatrix} m_0\text{: } y := y + 1 \\ m_1\text{:} \end{bmatrix}$$

linear invariant:

$$\varphi\text{:} \quad a \cdot y \ + \ b_{\ell_0} \cdot at\_\ell_0 \ + \ b_{\ell_1} \cdot at\_\ell_1 \ + \ b_{m_0} \cdot at\_m_0 \ + \ b_{m_1} \cdot at\_m_1 \ = \ K$$

(I) $a \cdot 0 + b_{\ell_0} + b_{m_0} = K$
(initial value of y is 0)

(T) $a \cdot 1 - b_{\ell_0} + b_{\ell_1} = 0$ (for ℓ_0)
$a \cdot 1 - b_{m_0} + b_{m_1} = 0$ (for m_0)

`Example:` Program DOUBLE (Con'd)

Possible solutions (basis for all solutions)

	a	b_{ℓ_0}	b_{ℓ_1}	b_{m_0}	b_{m_1}	K
S_1	0	1	1	0	0	1
S_2	0	0	0	1	1	1
S_3	1	1	0	1	0	2

Corresponding invariants

φ_1: $at\_\ell_0 + at\_\ell_1 = 1$ (control invariant)

φ_2: $at\_m_0 + at\_m_1 = 1$ (control invariant)

φ_3: $y + at\_\ell_0 + at\_m_0 = 2$

Linear Invariants for Cyclic Programs

Program $\ell_0^1\colon\ S_1 \parallel \ldots \parallel \ell_0^j\colon\ S_j \parallel \ldots \parallel \ell_0^m\colon\ S_m$
where S_j is of the form

$$\ell_0^j\colon\ \textbf{loop forever do}\ \underbrace{\ell_1^j, \ell_2^j, \ldots, \ell_k^j}_{\text{cycle } C}$$

Define

$$\Delta(y, C) = \Delta(y, \tau_1) + \cdots + \Delta(y, \tau_k)$$

For these programs construction of the linear invariants can be done in three phases:

1. Compute $\boldsymbol{a_i}$'s
2. Compute $\boldsymbol{b_\ell}$'s
3. Compute K

Phase 1: Bodies

For cycle $\underbrace{\ell_1, \ell_2, \ldots, \ell_k}_{C}$

$$
\begin{aligned}
\sum_{i=1}^{r} a_i \cdot \Delta(y_i, \tau_{\ell_1}) - b_{\ell_1} + b_{\ell_2} &= 0 \\
\sum_{i=1}^{r} a_i \cdot \Delta(y_i, \tau_{\ell_2}) - b_{\ell_2} + b_{\ell_3} &= 0 \\
&\vdots \\
\sum_{i=1}^{r} a_i \cdot \Delta(y_i, \tau_{\ell_k}) + b_{\ell_1} - b_{\ell_k} &= 0
\end{aligned}
$$

$$\sum_{i=1}^{r} a_i \cdot (\Delta(y_i, \tau_{\ell_1}) + \ \ldots \ + \Delta(y_i, \tau_{\ell_k})) = 0$$

Thus,

$$\boxed{\sum_{i=1}^{r} a_i \cdot \Delta(y_i, C) \ = \ 0}$$

Phase 2: Compensation Expressions

$$\boxed{b_{\ell_0} = 0}$$

For $\tau\colon \ell_j \to \ell_k$ where $j < k$

$$\sum_{i=1}^{r} a_i \cdot \Delta(y_i, \tau) \; - \; b_{\ell_j} \; + \; b_{\ell_k} \; = \; 0$$

Assume that for all $j < k$, b_{ℓ_j} is known.
Compute b_{ℓ_k} from

$$\boxed{b_{\ell_k} \; = \; b_{\ell_j} \; - \; \sum_{i=1}^{r} a_i \cdot \Delta(y_i, \tau)}$$

(independently for each cycle)

Phase 3: Right constants

$$\boxed{K = \sum_{i=1}^{r} a_i \cdot y_i^0}$$

Note: This set of equations has the same solutions as the equations (T) + (I) except for solutions of the form

$$at\_\ell_1 + \cdots + at\_\ell_k = 1$$

which are produced by (T) + (I), but not by this set.

`Example:` Program PROD-CON-SV (Fig 2.23)
Producer-Consumer with
shared variables

- semaphores r, ne, nf:

  ne – counts # of empty slots in list $\boldsymbol{b}$

  initially $ne = N$

  nf – counts # of full slots in $\boldsymbol{b}$

  initially $nf = 0$

  r – ensures that the shared variable $\boldsymbol{b}$ is handled exclusively by *Prod* or *Cons*

- linear variables: $r, ne, nf, |b|$

Program PROD-CONS-SV (Fig. 2.23)

$$
\begin{array}{ll}
\textbf{local}\ r,\ ne,\ nf & : \textbf{integer where}\ r = 1,\ ne = N,\ nf = 0 \\
\phantom{\textbf{local}}\ b & : \textbf{list of integer where}\ b = \Lambda
\end{array}
$$

$$
Prod :: \left[\begin{array}{l}
\textbf{local}\ x\text{: }\textbf{integer} \\
\ell_0\text{: }\textbf{loop forever do} \\
\qquad \left[\begin{array}{l}
\ell_1\text{: }\textbf{produce}\ x \\
\ell_2\text{: }\textbf{request}\ ne \\
\ell_3\text{: }\textbf{request}\ r \\
\ell_4\text{: }b := b \bullet x \\
\ell_5\text{: }\textbf{release}\ r \\
\ell_6\text{: }\textbf{release}\ nf
\end{array}\right]
\end{array}\right]
$$

$||$

$$
Cons :: \left[\begin{array}{l}
\textbf{local}\ y\text{: }\textbf{integer} \\
m_0\text{: }\textbf{loop forever do} \\
\qquad \left[\begin{array}{l}
m_1\text{: }\textbf{request}\ nf \\
m_2\text{: }\textbf{request}\ r \\
m_3\text{: }(y,\ b) := \big(hd(b),\ tl(b)\big) \\
m_4\text{: }\textbf{release}\ r \\
m_5\text{: }\textbf{release}\ ne \\
m_6\text{: }\textbf{consume}\ y
\end{array}\right]
\end{array}\right]
$$

Properties we want to prove:

$$\Box\, \underbrace{\neg(at\_\ell_4 \;\wedge\; at\_m_3)}_{\psi_1}$$

$$\Box\, \underbrace{at\_\ell_4 \;\rightarrow\; |b| < N}_{\psi_2}$$

$$\Box\, \underbrace{at\_m_3 \;\rightarrow\; |b| > 0}_{\psi_3}$$

Bottom-up invariants:

$$\underbrace{r \geq 0}_{\varphi_0} \;\wedge\; \underbrace{ne \geq 0}_{\varphi_1} \;\wedge\; \underbrace{nf \geq 0}_{\varphi_2} \;\wedge\; \underbrace{|b| \geq 0}_{\varphi_3}$$

<u>Bodies:</u>

Increments along each cycle:

	Prod	Cons
r	0	0
ne	-1	1
nf	1	-1
$\|b\|$	1	-1

For each cycle: $\sum_{i=1}^{r} a_i \cdot \Delta(y_i, C) \; = \; 0$

Therefore

Prod: $-a_e + a_f + a_b = 0$

Cons: $a_e - a_f - a_b = 0$

<u>Solutions</u>		<u>Bodies</u>
1. $a_r = 1,$	$a_e = a_f = a_b = 0$	B_1: r
2. $a_e = a_f = 1,$	$a_r = a_b = 0$	B_2: $ne + nf$
3. $a_e = a_b = 1,$	$a_r = a_f = 0$	B_3: $ne + \|b\|$

compensation expressions

coefficients of $b_{\ell_1}, \ldots, b_{m_6}$
corresponding to bodies
B_1: r
B_2: $ne + nf$
B_3: $ne + |b|$

– Prod –				– Cons –			
	B_1	B_2	B_3		B_1	B_2	B_3
b_{ℓ_1}	0	0	0	b_{m_1}	0	0	0
b_{ℓ_2}	0	0	0	b_{m_2}	0	1	0
b_{ℓ_3}	0	1	1	b_{m_3}	1	1	0
b_{ℓ_4}	1	1	1	b_{m_4}	1	1	1
b_{ℓ_5}	1	1	0	b_{m_5}	0	1	1
b_{ℓ_6}	0	1	0	b_{m_6}	0	0	0

Right constants

$b_{\ell_0} = b_{m_0} = 0$
Initial values

$$r = 1,\ ne = N,\ nf = 0,\ |b| = 0$$

$$K_1 \;=\; 1 \cdot \underbrace{1}_{r} \;=\; 1$$

$$K_2 \;=\; 1 \cdot \underbrace{N}_{ne} + 1 \cdot \underbrace{0}_{nf} \;=\; N$$

$$K_3 \;=\; 1 \cdot \underbrace{N}_{ne} + 1 \cdot \underbrace{0}_{|b|} \;=\; N$$

The resulting invariants

$$\boxed{\begin{array}{ll} \alpha_1: & r + at\_\ell_{4,5} + at\_m_{3,4} = 1 \\ \alpha_2: & ne + nf + at\_\ell_{3..6} + at\_m_{2..5} = N \\ \alpha_3: & ne + |b| + at\_\ell_{3,4} + at\_m_{4,5} = N \end{array}}$$

No need to check invariance!

These invariants imply the properties we wanted to prove:

$$\psi_1: \quad \underbrace{r + at\_\ell_{4,5} + at\_m_{3,4} = 1}_{\alpha_1} \wedge \underbrace{r \geq 0}_{\varphi_0}$$

$$\rightarrow \underbrace{\neg(at\_\ell_4 \wedge at\_m_4)}_{\psi_1}$$

$$\psi_2: \quad \underbrace{ne + |b| + at\_\ell_{3,4} + at\_m_{4,5} = N}_{\alpha_3} \wedge \underbrace{ne \geq 0}_{\varphi_1}$$

$$\rightarrow \underbrace{at\_\ell_4 \rightarrow |b| < N}_{\psi_2}$$

Since $at\_\ell_4 \rightarrow at\_\ell_{3,4} = 1$
and $ne \geq 0$, $at\_\ell_{3,4} = 1$, $at\_m_{4,5} \geq 0$ implies $|b| < N$

$$\psi_3: \quad \underbrace{ne + nf + at\_\ell_{3..6} + at\_m_{2..5} = N}_{\alpha_2} \;\wedge$$

$$\underbrace{ne + |b| + at\_\ell_{3,4} + at\_m_{4,5} = N}_{\alpha_3} \;\wedge$$

$$\underbrace{nf \geq 0}_{\varphi_2}$$

$$\rightarrow \; \underbrace{at\_m_3 \rightarrow |b| > 0}_{\psi_3}$$

Suppose at_m_3:

φ_2: $\quad ne + nf + at\_\ell_{3..6} + 1 = N$

φ_3: $\quad ne + |b| + at\_\ell_{3,4} + 0 = N$

Since $\varphi_2 - \varphi_3$ yields
$nf - |b| + at\_\ell_{3..6} - at\_\ell_{3,4} + 1 = 0$

Thus
$$|b| = \underbrace{nf}_{\geq 0} + \underbrace{(at\_\ell_{3..6} - at\_\ell_{3,4})}_{\geq 0} + 1 > 0$$