CS256/Spring 2008 — Lecture #14

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Satisfiability over a finite-state program

P-validity problem (of φ)

Given a finite-state program P
and formula φ,

is φ P-valid?

i.e. do all P-computations satisfy φ?

P-satisfiability problem (of φ)

Given a finite-state program P
and formula φ

is φ P-satisfiable?

i.e., does there exist a P-computation which satisfies φ?

To determine whether φ is P-valid,
it suffices to apply an algorithm for
deciding if there is a P-computation
that satisfies $\neg\varphi$.

The Idea

To check P-satisfiability of φ,
we combine the tableau T_φ and the
transition graph G_P into one product graph,
called the behavior graph $\mathcal{B}_{(P,\varphi)}$,
and search for paths

$$(s_0, A_0),\ (s_1, A_1),\ (s_2, A_2),\ \ldots$$

that satisfy the two requirements:

- $\sigma \vDash \varphi$:
  there exists a fulfilling path
  $\pi:\ A_0, A_1, \ldots$
  in the tableau T_φ such that $\varphi \in A_0$.
- σ is a P-computation:
  there exists a fair path
  $\sigma:\ s_0, s_1, \ldots$
  in the transition graph G_P.

State transition graph G_P: Construction

- Place as nodes in G_P all initial states s ($s \Vdash \Theta$)
- Repeat

  for some $s \in G_P,\ \tau \in \mathcal{T}$,
  add all its τ-successors s' to G_P
  if not already there,
  and add edges between s and s'.

  Until no new states or edges can be added.

If this procedure terminates, the system is finite-state.

Example: Program mux-pet1 (Fig. 3.4)

(Peterson's Algorithm for mutual exclusion)

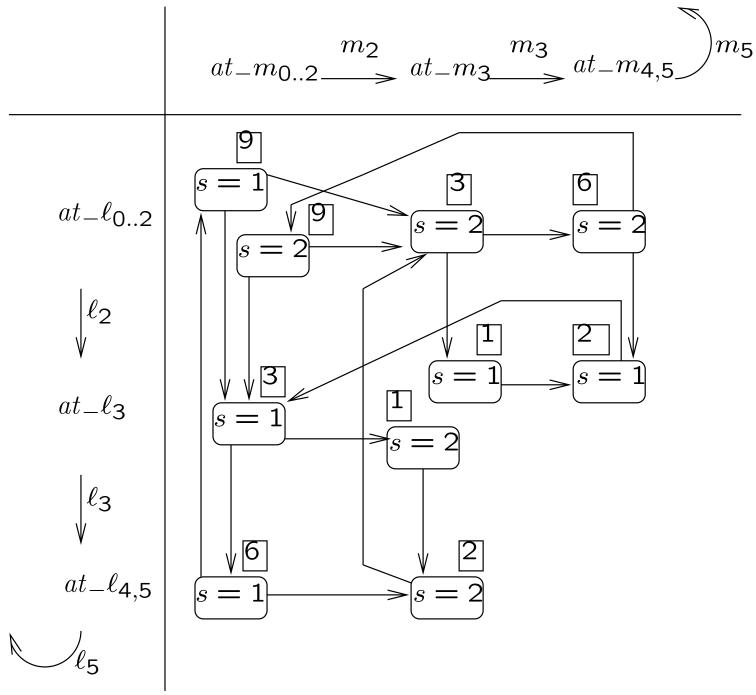
local y_1, y_2: **boolean** **where** $y_1 = \text{F}, y_2 = \text{F}$
s : **integer** **where** $s = 1$

$$P_1 :: \quad \ell_0 : \textbf{loop forever do} \left[\begin{array}{ll} \ell_1 : & \textbf{noncritical} \\ \ell_2 : & (y_1,\ s) := (\text{T},\ 1) \\ \ell_3 : & \textbf{await } (\neg y_2) \vee (s \neq 1) \\ \ell_4 : & \textbf{critical} \\ \ell_5 : & y_1 := \text{F} \end{array} \right]$$

$||$

$$P_2 :: \quad m_0 : \textbf{loop forever do} \left[\begin{array}{ll} m_1 : & \textbf{noncritical} \\ m_2 : & (y_2,\ s) := (\text{T},\ 2) \\ m_3 : & \textbf{await } (\neg y_1) \vee (s \neq 2) \\ m_4 : & \textbf{critical} \\ m_5 : & y_2 := \text{F} \end{array} \right]$$

Abstract state-transition graph for MUX-PET1



We use $y_1 \Leftrightarrow at\_\ell_{3..5}$
$y_2 \Leftrightarrow at\_m_{3..5}$

Some states have been lumped together:
a super-state labeled by $\boxed{i}$ represents i states

MUX-PET1 has 42 reachable states.

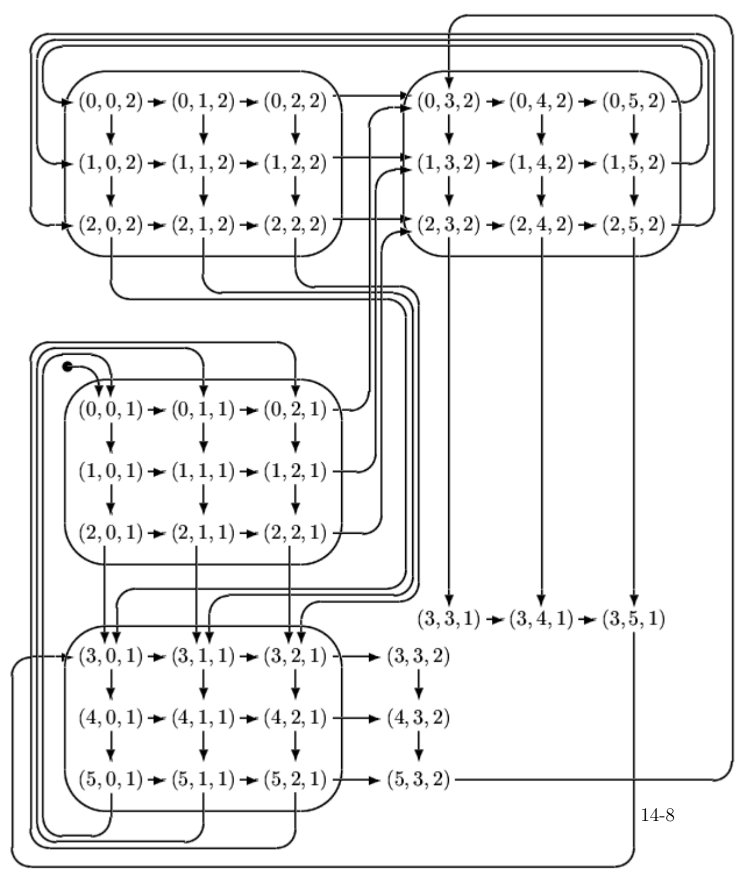
Based on this graph it is straightforward to check the properties

ψ_1 : $\Box\neg(at\_\ell_4 \wedge at\_m_4)$

ψ_2 : $\Box(at\_\ell_3 \wedge \neg at\_m_3 \rightarrow s = 1)$

ψ_3 : $\Box(at\_m_3 \wedge \neg at\_\ell_3 \rightarrow s = 2)$

MUX-PET1 Full state-transition graph (l_i, m_j, s)



Definitions

- For atom A, $state(A)$ is the conjunction of all state formulas in A
  (by R_{sat}, $state(A)$ must be satisfiable)

- For $A \in T_\varphi$,
  $\delta(A)$ denotes the set of successors of A in T_φ

- atom A is consistent with state s
  if $s \Vdash state(A)$,

  i.e. s satisfies all state formulas in A.

- ϑ: $A_0, A_1, \ldots$ path in T_φ
  σ: $s_0, s_1, \ldots$ computation of P

  ϑ is a trail of T_φ over σ if
  A_j is consistent with s_j, for all $j \geq 0$

Behavior Graph

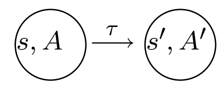
For finite-state program P and formula φ,
we construct the (P, φ)-behavior graph

$$\mathcal{B}_{(P,\varphi)} \quad \approx \quad G_P \times T_\varphi^- \text{ (pruned)}$$

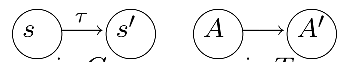
such that

- nodes are labeled by (s, A)
  where s is a state from G_P and
  A is an atom from T_φ consistent with s.

- edges
  There is an edge



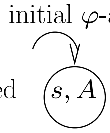
  if and only if $s' \in \tau(s)$ and $A' \in \delta(A)$



in G_P in T_φ

- initial φ-node $\quad (s, A)$

  if s is an initial state $(s \Vdash \Theta)$
  and A is an initial φ-atom $(\varphi \in A)$

  It is marked 

Algorithm behavior-graph
(constructing $\mathcal{B}_{(P,\varphi)}$)

- Place in $\mathcal{B}$ all initial φ-nodes (s, A)
  (s initial state of P,
  A initial φ-atom in T_φ^-
  A consistent with s)

- Repeat until no new nodes or new edges can be added:

  Let (s, A) be a node in $\mathcal{B}$
  $\tau \in \mathcal{T}$ a transition
  (s', A') a pair s.t.
  s' is a τ-successor of s
  $A' \in \delta(A)$ in pruned T_φ^-
  A' consistent with s'

  - Add (s', A') to $\mathcal{B}$, if not already there
  - Draw a τ-edge from (s, A) to (s', A'), if not already there

`Example:` Given FTS LOOP

$\Theta: \; x = 0$

$\mathcal{T} = \{\tau, \tau_I\}$

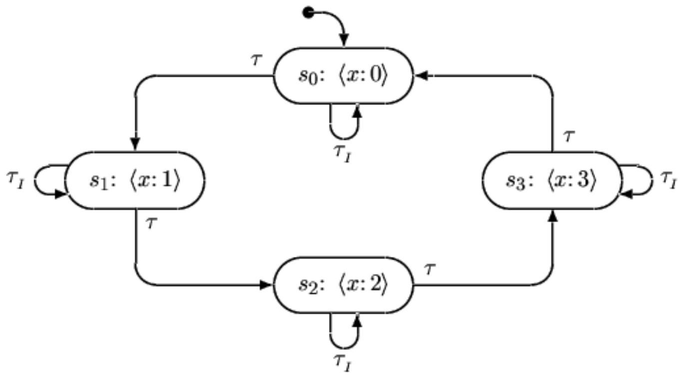
with τ_I (idling)
τ where ρ_τ: $x' = (x+1) mod 4$

$\mathcal{J}$: $\{\tau\}$

Check P-satisfiability of $\boxed{\psi_3\text{: } \Diamond \Box (x \neq 3)}$

state-transition graph G_{LOOP} (Fig 5.9)
pruned $T_{\psi_3}^-$ (Fig 5.8)
Behavior graph $\mathcal{B}_{(\text{LOOP},\psi_3)}$ (Fig 5.10)

Fig. 5.9. State-transition graph G_{LOOP}

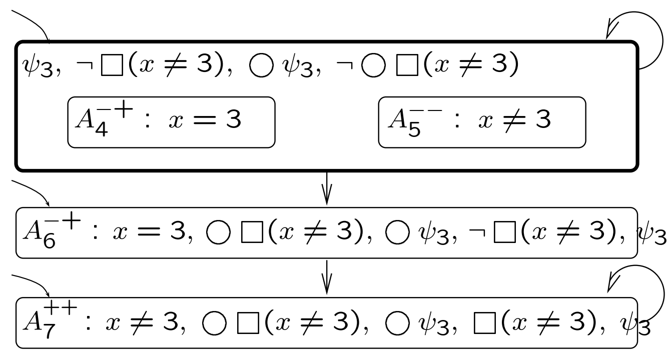


<u>Pruned tableau $T^{-}_{\psi_3}$</u> (Fig. 5.8)

Eliminating
- MSCS's not reachable from an initial ψ_3-atom and
- non-fulfilling terminal MSCS's

Promising formulas:

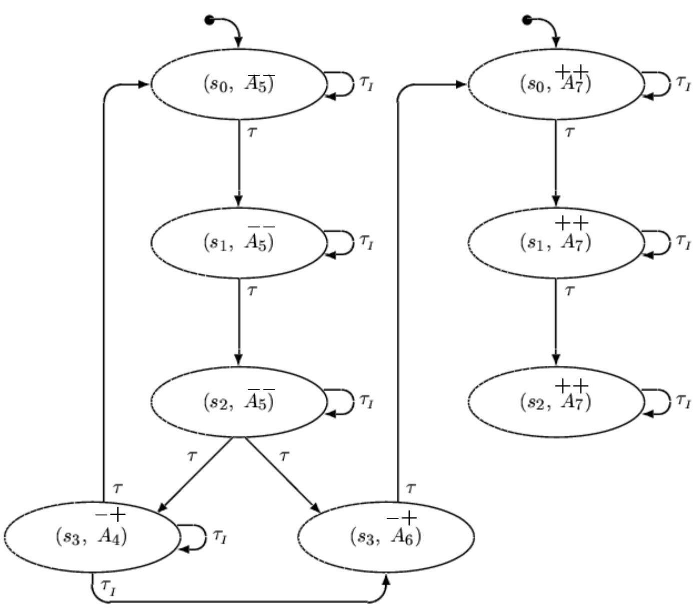
$$\Diamond\,\Box(x \neq 3) \quad \text{promising} \quad \Box(x \neq 3)$$
$$\neg\,\Box(x \neq 3) \quad \text{promising} \quad (x = 3)$$



Two non-transient MSCS's:

$\{A_4^{-+}, A_5^{--}\}$	not fulfilling
$\{A_7^{++}\}$	fulfilling

Behavior graph $\mathcal{B}_{(\text{LOOP},\psi_3)}$ (Fig 5.10)



<u>`Example:`</u> Given FTS ONE:

Θ: $x = 0$

$\mathcal{T}$: $\{\tau_1, \tau_2, \tau_3, \tau_4, \tau_I\}$

with $\rho_{\tau_1}:\ x = 0 \wedge x' = 1$

$\rho_{\tau_2}:\ x = 1 \wedge x' = 0$

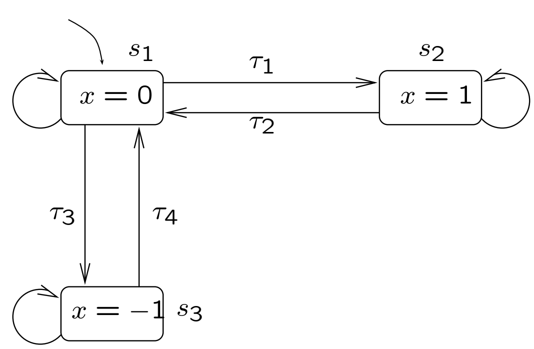
$\rho_{\tau_3}:\ x = 0 \wedge x' = -1$

$\rho_{\tau_4}:\ x = -1 \wedge x' = 0$

$\mathcal{J}$: $\emptyset$

$\mathcal{C}$: $\{\tau_1, \tau_3\}$

Transition graph G_{ONE}



We want to know whether

$$\boxed{\varphi:\ \square\diamond(x=1)}$$

is valid over ONE.

Check P-satisfiability of

$$\boxed{\neg\varphi:\ \underbrace{\diamond\square(x\neq 1)}_{\psi}}$$

$\Phi_\psi^+:\ \{\psi,\ \bigcirc\psi,\ \square(x\neq 1),\ \bigcirc\square(x\neq 1),\ x=1\}$

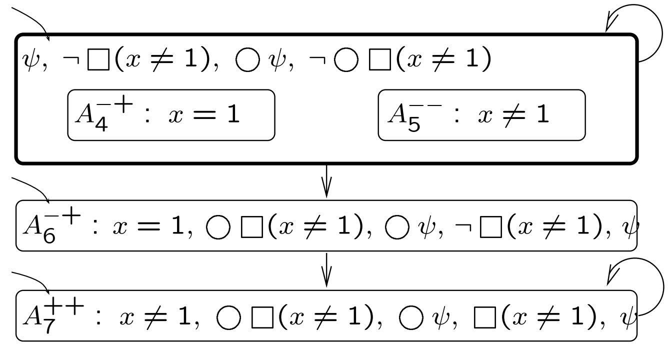
basic formulas: $\{\bigcirc\psi,\ \bigcirc\square(x\neq 1),\ x=1\}$

Promising formulas:

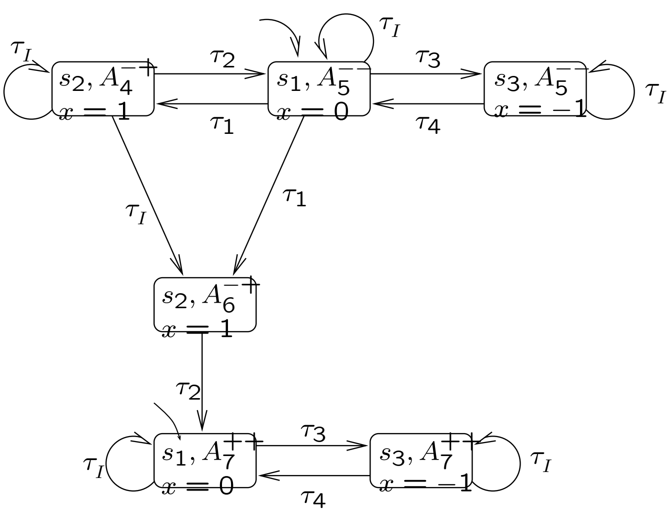
$\psi_1:\psi=\diamond\square(x\neq 1)$ promising $r_1:\ \square(x\neq 1)$

$\psi_2:\ \neg\square(x\neq 1)$ promising $r_2:\ x=1$

Pruned tableau T_ψ^-



Behavior graph $\mathcal{B}_{(\text{ONE},\diamond\square(x\neq 1))}$



Two non-transient MSCS's:

$\{(s_2,A_4^{-+}),\ (s_1,A_5^{--}),\ (s_3,A_5^{--})\}$: not fulfilling,

$\{(s_1,A_7^{++}),(s_3,A_7^{++})\}$: fulfilling

Paths of $\mathcal{B}_{(P,\varphi)}$

`Claim 5.9` (paths of $\mathcal{B}_{(P,\varphi)}$)

The infinite sequence

$$\pi:\ \underbrace{(s_0,A_0)}_{\varphi\text{-initial}},\ (s_1,A_1),\ \ldots$$

is a path in $\mathcal{B}_{(P,\varphi)}$

iff

$\sigma_\pi:\ s_0,s_1,\ldots$ is a run of P
(i.e. computation of P less fairness)

$\vartheta_\pi:\ A_0,A_1,\ldots$ is a trail of T_φ over σ_π
(i.e. A_j consistent with s_j, for all $j\geq 0$)

`Example:` In $\mathcal{B}_{(\text{LOOP},\psi_3)}$ (Fig. 5.10)

$\pi:\ \big((s_0,A_5),\ (s_1,A_5),\ (s_2,A_5),\ (s_3,A_4)\big)^\omega$

induces

$\sigma_\pi:\ (s_0,s_1,s_2,s_3)^\omega$ run of LOOP

$\vartheta_\pi:\ (A_5,A_5,A_5,A_4)^\omega$ trail of T_{ψ_3} over σ_π

Proposition 5.10 (P-satisfiability by path)

P has a computation satisfying φ
iff
there is an infinite φ-initialized path π
in $\mathcal{B}_{(P,\varphi)}$ s.t.

σ_π is a P-computation (fair run of P)

ϑ is a fulfilling trail over σ_π

Searching for "good" paths in $\mathcal{B}_{(P,\varphi)}$
— not practical.

Definitions

For behavior graph $\mathcal{B}_{(P,\varphi)}$

- node (s', A') is a τ-successor of (s, A)
  if $\mathcal{B}_{(P,\varphi)}$ contains τ-edge connecting (s, A) to (s', A')

- transition τ is enabled on node (s, A)
  if τ is enabled on state s

Definitions (Con't)

For SCS $S \subseteq \mathcal{B}_{(P,\varphi)}$:

- Transition τ is taken in S if there exists two nodes $(s, A), (s', A') \in S$ s.t.
  (s', A') is a τ-successor of (s, A)

- S is $\left\{\begin{matrix}\text{just}\\ \text{compassionate}\end{matrix}\right\}$ if every $\left\{\begin{matrix}\text{just}\\ \text{compassionate}\end{matrix}\right\}$ transition $\tau \left\{\begin{matrix}\in \mathcal{J}\\ \in \mathcal{C}\end{matrix}\right\}$ is either taken in S or is disabled on $\left\{\begin{matrix}\text{some node}\\ \text{all nodes}\end{matrix}\right\}$ in S

- S is fair if it is both just and compassionate

- S is fulfilling if every promising formula $\psi \in \Phi_\psi$ is fulfilled by some atom A, s.t.
  $(s, A) \in S$ for some state s

- S is adequate if it is fair and fulfilling

Adequate SCS's

Proposition 5.11 (adequate SCS and satisfiability)

Given a finite-state program P and temporal formula φ.
φ is P-satisfiable
iff

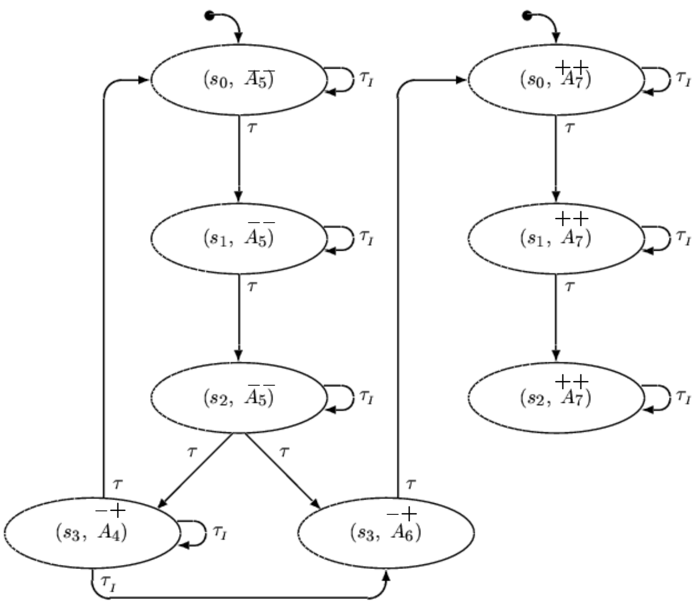
$\mathcal{B}_{(P,\varphi)}$ has an adequate SCS

Example: Consider LOOP and

$$\psi_3: \Diamond \Box (x \neq 3)$$

Is ψ_3 LOOP-satisfiable?
Check the SCS's in $\mathcal{B}_{(\text{LOOP},\psi_3)}$ (Fig. 5.10)

Behavior graph $\mathcal{B}_{(\text{LOOP},\psi_3)}$ (Fig 5.10)



Example (Con't)

- $\{\ (s_0, A_5^{--}),\ (s_1, A_5^{--}),\ (s_2, A_5^{--}),\ (s_3, A_4^{-+})\ \}$
  is fair but not fulfilling

- $\{\ (s_0, A_7^{++})\},\ \{(s_1, A_7^{++})\},\ \{(s_2, A_7^{++})\}$
  each is fulfilling but not fair
  Not just with respect to transition τ

- $\{(s_3, A_6^{-+})\}$
  is neither fair (unjust toward τ) nor fulfilling (being transient)

No adequate subgraphs in $\mathcal{B}_{(\text{LOOP},\psi_3)}$

Therefore, by **proposition 5.11**, LOOP has no computation that satisfies ψ_3: $\diamond\square(x \neq 3)$

Example: Consider LOOP and

φ_3: $\square\diamond(x = 3)$

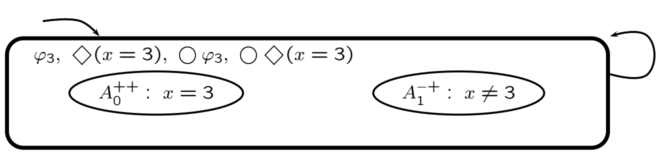
Is φ_3 LOOP-satisfiable?

Promising formulas :

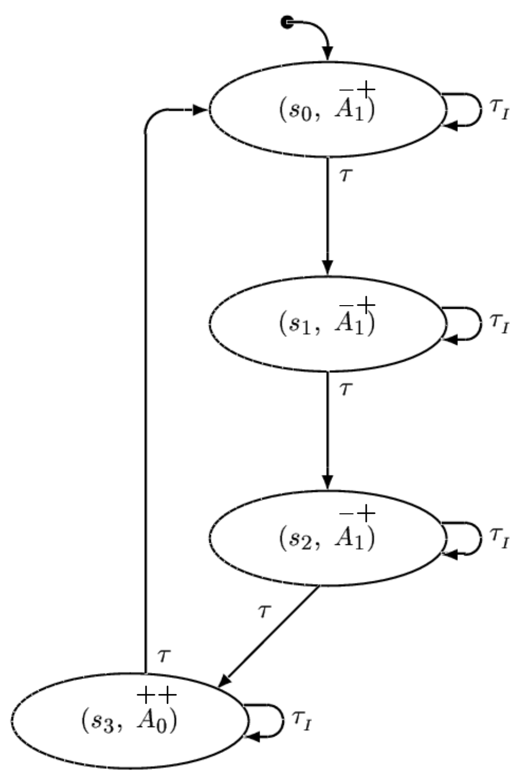
$\diamond(x = 3)$ promising $(x = 3)$

$\neg\square\diamond(x = 3)$ promising $\neg\diamond(x = 3)$

Pruned tableau T_{φ_3} (Fig. 5.6)



Behavior graph $\mathcal{B}_{(\text{LOOP},\varphi_3)}$ (Fig. 5.11)



$S = \{ (s_0, A_1^{-+}),\ (s_1, A_1^{-+}),\ (s_2, A_1^{-+}),\ (s_3, A_0^{++}) \}$

is an adequate subgraph:

fair $\qquad$ (τ taken in S)

fulfilling

Therefore, by `proposition 5.11`, program LOOP has a computation satisfying φ_3: $\Box \Diamond (x = 3)$

The periodic computation σ: $(x{:}\,0, x{:}\,1, x{:}\,2, x{:}\,3)^\omega$ satisfies φ_3

From Atom Tableau T_φ to ω-Automaton $\mathcal{A}_\varphi$

For temporal formula φ, construct the ω-automaton

$$\mathcal{A}_\varphi :\ \langle \underbrace{N,\ N_0,\ E,}_{\substack{\text{Same as}\\ T_\varphi}}\ \mu,\ \mathcal{F} \rangle$$

where

- Node labeling μ:
  For node $n \in N$ labeled by atom A in T_φ,
  $$\mu(n) \ = \ state(A).$$

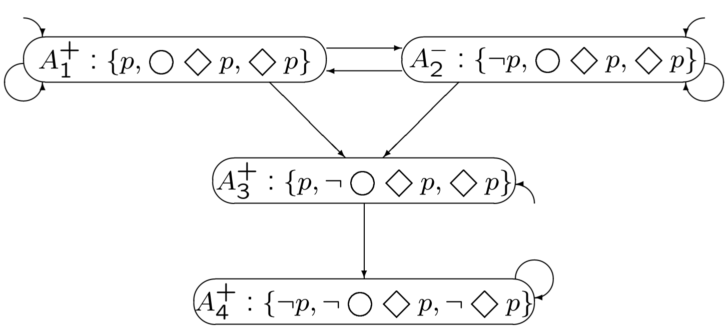
- Acceptance condition $\mathcal{F}$:

  Muller:
  $$\mathcal{F} \ = \ \{\text{SCS } S \mid S \text{ is fulfilling } \}$$

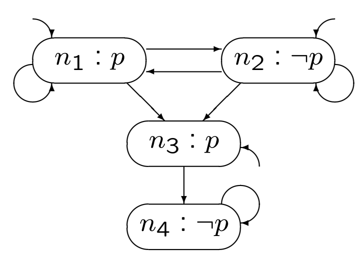
  Street:
  $$\mathcal{F} \ = \ \{(P_\psi, R_\psi) \mid \psi \in \Phi_\varphi \text{ promises } r\},$$
  where
  $$\begin{aligned} P_\psi &= \{\ A \mid \neg\psi \in A\ \} \\ R_\psi &= \{\ A \mid r \in A\ \} \end{aligned}$$

`Example`: $\varphi :\ \Diamond p$

Tableau T_φ:



`Example`: $\mathcal{A}_{\Diamond p}$ from $T_{\Diamond p}$



$$\begin{aligned} \mathcal{F}_M &= \{\{n_1\}, \{n_1, n_2\}, \{n_4\}\} \\ \mathcal{F}_S &= \{(P_{\Diamond p}, R_{\Diamond p})\} \\ &= \{(\{n_4\}, \{n_1, n_3\})\} \\ &\approx \{(\{n_4\}, \{n_1\})\} \end{aligned}$$

since no path can visit n_3 infinitely often

Abstraction

Abstraction = a method to verify infinite-state systems.

Idea:

	abstraction ↓	
Program P (infinite state)	$\longrightarrow$	Abstract program P^A (finite state)
Property φ	$\longrightarrow$	Abstract property φ^A
$P \vDash \varphi$?	$\longrightarrow$	$P^A \vDash \varphi^A$ ↓ model checking

We want to ensure that
if $P^A \vDash \varphi^A$ then $P \vDash \varphi$.

Abstraction (Cont'd)

How do we obtain such an abstraction function?

- 1) Abstract the domain to a finite-state one (*data abstraction*):
  For variables $\vec{x}$ ranging over domain D, find an abstract domain D^A and an abstraction function $\alpha : D \rightarrow D^A$.

- 2) From the data abstraction it is possible to compute an abstraction for the program and for the property such that
  if $P^A \vDash \varphi^A$ then $P \vDash \varphi$.

Example: Abstracting Bakery

Program MUX-BAK (infinite-state program)

$$P_1 :: \left[\begin{array}{l}\textbf{loop forever do}\\ \left[\begin{array}{l}\ell_0 : \textbf{noncritical}\\ \ell_1 : y_1 := y_2 + 1\\ \ell_2 : \textbf{await } y_2 = 0 \lor y_1 \le y_2\\ \ell_3 : \textbf{critical}\\ \ell_4 : y_1 := 0\end{array}\right]\end{array}\right]$$

$\|$

$$P_2 :: \left[\begin{array}{l}\textbf{loop forever do}\\ \left[\begin{array}{l}m_0 : \textbf{noncritical}\\ m_1 : y_2 := y_1 + 1\\ m_2 : \textbf{await } y_1 = 0 \lor y_2 < y_1\\ m_3 : \textbf{critical}\\ m_4 : y_2 := 0\end{array}\right]\end{array}\right]$$

Abstract domain: the boolean algebra over
$B = \{b_1, b_2, b_3 : \textbf{boolean}\}$,
with $b_1 : \; y_1 = 0$
$b_2 : \; y_2 = 0$
$b_3 : \; y_1 \le y_2$

Example: Abstracting Bakery (Cont'd)

Program MUX-BAK-ABSTR (finite-state program)

$$P_1 :: \left[\begin{array}{l}\textbf{loop forever do}\\ \left[\begin{array}{l}\ell_0 : \textbf{noncritical}\\ \ell_1 : (b_1, b_3) := (\mathit{false}, \mathit{false})\\ \ell_2 : \textbf{await } b_2 \lor b_3\\ \ell_3 : \textbf{critical}\\ \ell_4 : (b_1, b_3) := (\mathit{true}, \mathit{true})\end{array}\right]\end{array}\right]$$

$\|$

$$P_2 :: \left[\begin{array}{l}\textbf{loop forever do}\\ \left[\begin{array}{l}m_0 : \textbf{noncritical}\\ m_1 : (b_2, b_3) := (\mathit{false}, \mathit{true})\\ m_2 : \textbf{await } b_1 \lor \neg b_3\\ m_3 : \textbf{critical}\\ m_4 : (b_2, b_3) := (\mathit{true}, b_1)\end{array}\right]\end{array}\right]$$

This program can now be checked for mutual exclusion, bounded overtaking, response.

Show MUX-BAK-ABSTR $\vDash \Box\, \neg(at\_\ell_3 \land at\_m_3)$. Then it follows that MUX-BAK $\vDash \Box\, \neg(at\_\ell_3 \land at\_m_3)$.