

Circuit: A Regent Application

CS315B
Lecture 6

Circuit

- Electrical simulation
- A graph
 - Wires are edges
 - Nodes are places where wires meet

Circuit.rg

- Iterative simulation with three phases:
 - calculate_new_currents
 - distribute_charge
 - update_voltages
- New features
 - Structs
 - Permissions on multiple fields
 - wait_for(...)
 - __demand(...)

Circuit_dep_par.rg

- New features
 - Pointers to region unions
 - Reduce privilege
 - __demand(__parallel)
 - Tracing

Circuit Dependent Partitioning

```
var pn_equal = partition(equal, rn, colors)
var pw_outgoing = preimage(rw, pn_equal, rw.in_ptr)
var pw_incoming = preimage(rw, pn_equal, rw.out_ptr)
var pw_crossing_out = pw_outgoing - pw_incoming
var pw_crossing_in = pw_incoming - pw_outgoing
var pn_shared_in = image(rn, pw_crossing_in, rw.out_ptr)
var pn_shared_out = image(rn, pw_crossing_out, rw.in_ptr)
var pn_private = (pn_equal - pn_shared_in) - pn_shared_out
var pn_shared = pn_equal - pn_private
var pn_ghost = image(rn, pw_crossing_out, rw.out_ptr)
```

Mapping

Mapping

- Mapping is the process of assigning resources to Regent/Legion programs
- Conceptually
 - Assign a processor to each task
 - The task will execute in its entirety on that processor
 - Assign a memory to each region argument
- Can also control other things
 - But these are the most important

The Legion Mapping API

- A *mapper* implements the Legion mapping API
 - A set of C++ callbacks
- Bishop is a high-level language for mapping
 - Compiles to the mapping API
 - Not as general, but easier to learn and use

High-Level Overview

- An instance of the Legion runtime runs on every node
- When a task is launched the local runtime
 - Makes mapper calls to pick a processor for the task
 - Makes mapper calls to pick memories for the region arguments
 - ... and other mapper calls as well ...

New Concepts

- There are a number of concepts in mapping that don't exist in Regent
- Machine models
- Variants
- Physical Instances

Machine Model

- To pick concrete processors & memories, the runtime must know:
- How many processors/memories there are
 - And of what kinds
- And where the processors/memories are
 - At least relative to each other

Machine Model

- | | |
|--------------|----------------|
| • Processors | • Memories |
| - LOC | - GLOBAL |
| - TOC | - SYSTEM |
| - PROC_SET | - RDMA |
| - UTILITY | - FRAME_BUFFER |
| - IO | - ZERO_COPY |
| | - DISK |
| | - HDF5 |

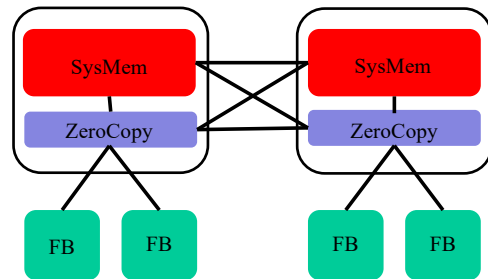
Affinities

- Processor → Memory
 - Which memories are attached to a processor
- Memory → Memory
 - Which memories have channels between them
- Memory → Processor
 - All processors attached to a memory
- Affinities are provided as a list of $(proc, mem)$ and (mem, mem) pairs

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An Example Machine Model (Simplified)



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Task Variants

- A task can have multiple *variants*
 - Different implementations of the same task
 - Multiple variants can be registered with the runtime
- Examples
 - A variant for LOC
 - Another variant for TOC
 - Variants for different data layouts

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Physical Instances

- A *region* is a logical name for data
- A *physical instance* is a copy of that data
 - For some set of fields
- There can be 0, 1 or many physical instances of a specific field of a region at any time

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Physical Instances

- Can be *valid* or *invalid*
 - Is the data current or not?
- Live in a specific memory
- Have a specific layout
 - Column major, row major, blocked, struct-of-arrays, array-of-structs, ...
- Are allocated explicitly by the mapper
- Are deallocated by the runtime
 - Garbage collected

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Index Launches

- A normal task call launches a single task
- An *index task call* launches a set of tasks
 - One for each point in a supplied index space
- Index launches are more efficient than launching many tasks individually
 - Regent automatically transforms loops of single task launches into index task launches

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Example

```
for x in prt.colors do  
  task(prt[x])
```

becomes

```
index_launch(task,prt,prt.colors)
```

(if there are no dependencies)

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A Mapper

- The circuit custom mapper, `circuit_bishop.rg`

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Miscellaneous Mapping Topics ...

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Controlling Processor Choice in Regent

- Place immediately before a task declaration
 - `__demand(__cuda)`
- Causes both CPU and GPU task variants to be produced
- And the default mapper always prefers to pick a GPU variant if possible

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Layout Constraints

- Tasks can have layout constraints on physical instances
 - "This task requires data in row major order"
- Constraints are just that
 - Don't specify an exact layout
 - Multiple instances may satisfy the constraints

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Reduction Instances

- A *reduction instance* is a special instance used for reductions

```
fill(R', 0)
for i in R.indices do
  R'[i] += val1
  R'[i] += val2
```
- Pattern

```
for i in R do
  i.field += val1
  i.field += val2

... later ...

R += R'
```

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Virtual Mappings

- It is also possible for a mapper to map a region to *no* instance
 - If the task does not use the region itself
 - E.g., only passes it to subtasks
- This is a *virtual mapping*

Summary

- Mapping
 - Selects processors for tasks
 - Selects memories for physical instances
 - Satisfying region requirements of tasks
- Choices
 - Use Bishop - easy to write a custom mapper
 - Default mapper does reasonable things
 - Overriding methods in the default mapper provides more options than Bishop mappers