

## Circuit: A Regent Application

CS315B  
Lecture 6

## Circuit

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- Electrical simulation
- A graph
  - Wires are edges
  - Nodes are places where wires meet

## Circuit.rg

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- Iterative simulation with three phases:
  - calculate\_new\_currents
  - distribute\_charge
  - update\_voltages
- New features
  - Structs
  - Permissions on multiple fields
  - wait\_for(...)
  - \_\_demand(...)

## Circuit\_dep\_par.rg

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- New features
  - Pointers to region unions
  - Reduce privilege
  - \_\_demand(\_\_parallel)
  - Tracing

## Circuit Dependent Partitioning

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```
var pn_equal = partition(equal, rn, colors)
var pw_outgoing = preimage(rw, pn_equal, rw.in_ptr)
var pw_incoming = preimage(rw, pn_equal, rw.out_ptr)
var pw_crossing_out = pw_outgoing - pw_incoming
var pw_crossing_in = pw_incoming - pw_outgoing
var pn_shared_in = image(rn, pw_crossing_in, rw.out_ptr)
var pn_shared_out = image(rn, pw_crossing_out, rw.in_ptr)
var pn_private = (pn_equal - pn_shared_in) - pn_shared_out
var pn_shared = pn_equal - pn_private
var pn_ghost = image(rn, pw_crossing_out, rw.out_ptr)
```

## Mapping

## Mapping

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- Mapping is the process of assigning resources to Regent/Legion programs
- Conceptually
  - Assign a processor to each task
    - The task will execute in its entirety on that processor
  - Assign a memory to each region argument
- Can also control other things
  - But these are the most important

## The Legion Mapping API

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- A *mapper* implements the Legion mapping API
  - A set of C++ callbacks
- Bishop is a high-level language for mapping
  - Compiles to the mapping API
  - Not as general, but easier to learn and use

## High-Level Overview

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- An instance of the Legion runtime runs on every node
- When a task is launched the local runtime
  - Makes mapper calls to pick a processor for the task
  - Makes mapper calls to pick memories for the region arguments
  - ... and other mapper calls as well ...

## New Concepts

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- There are a number of concepts in mapping that don't exist in Regent
- Machine models
- Variants
- Physical Instances

## Machine Model

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- To pick concrete processors & memories, the runtime must know:
- How many processors/memories there are
  - And of what kinds
- And where the processors/memories are
  - At least relative to each other

## Machine Model

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- Processors
  - LOC
  - TOC
  - PROC\_SET
  - UTILITY
  - IO
- Memories
  - GLOBAL
  - SYSTEM
  - RDMA
  - FRAME\_BUFFER
  - ZERO\_COPY
  - DISK
  - HDF5

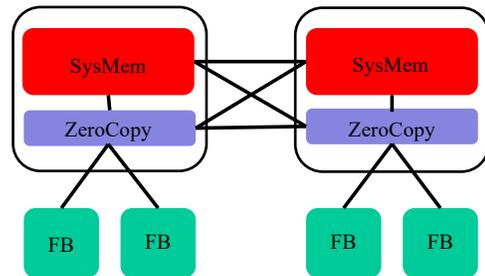
## Affinities

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- Processor → Memory
  - Which memories are attached to a processor
- Memory → Memory
  - Which memories have channels between them
- Memory → Processor
  - All processors attached to a memory
- Affinities are provided as a list of  $(proc, mem)$  and  $(mem, mem)$  pairs

## An Example Machine Model (Simplified)

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## Task Variants

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- A task can have multiple *variants*
  - Different implementations of the same task
  - Multiple variants can be registered with the runtime
- Examples
  - A variant for LOC
  - Another variant for TOC
  - Variants for different data layouts

## Physical Instances

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- A *region* is a logical name for data
- A *physical instance* is a copy of that data
  - For some set of fields
- There can be 0, 1 or many physical instances of a specific field of a region at any time

## Physical Instances

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- Can be *valid* or *invalid*
  - Is the data current or not?
- Live in a specific memory
- Have a specific layout
  - Column major, row major, blocked, struct-of-arrays, array-of-structs, ...
- Are allocated explicitly by the mapper
- Are deallocated by the runtime
  - Garbage collected

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## Index Launches

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- A normal task call launches a single task
- An *index task call* launches a set of tasks
  - One for each point in a supplied index space
- Index launches are more efficient than launching many tasks individually
  - Regent automatically transforms loops of single task launches into index task launches

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## Example

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```
for x in prt.colors do
  task(prt[x])
```

becomes

```
index_launch(task,prt,prt.colors)
```

(if there are no dependencies)

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## A Mapper

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- The circuit custom mapper, `circuit_bishop.rg`

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## Miscellaneous Mapping Topics ...

## Controlling Processor Choice in Regent

- Place immediately before a task declaration
  - `__demand(__cuda)`
- Causes both CPU and GPU task variants to be produced
- And the default mapper always prefers to pick a GPU variant if possible

## Layout Constraints

- Tasks can have layout constraints on physical instances
  - "This task requires data in row major order"
- Constraints are just that
  - Don't specify an exact layout
  - Multiple instances may satisfy the constraints

## Reduction Instances

- A *reduction instance* is a special instance used for reductions
  - `fill(R', 0)`
  - `for i in R.indices do`
    - `R'[i] += val1`
    - `R'[i] += val2`
- Pattern
  - `for i in R do`
    - `i.field += val1`
    - `i.field += val2`
  - ... later ...
  - `R += R'`

## Virtual Mappings

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- It is also possible for a mapper to map a region to *no* instance
  - If the task does not use the region itself
  - E.g., only passes it to subtasks
- This is a *virtual mapping*

## Summary

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- Mapping
  - Selects processors for tasks
  - Selects memories for physical instances
    - Satisfying region requirements of tasks
- Choices
  - Use Bishop - easy to write a custom mapper
  - Default mapper does reasonable things
    - Overriding methods in the default mapper provides more options than Bishop mappers