

Regent: More on Regions

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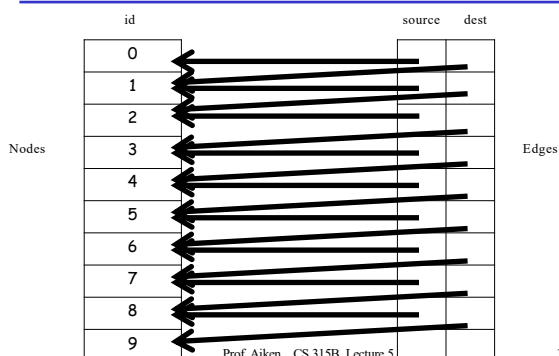
Regions Review

- A region is a (typed) collection
- Regions are the cross product of
 - An *index space*
 - A *field space*
- So far we've seen regions with N-dim index spaces
 - E.g., int1d, int2d, int3d

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Example 19



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Equal Partitioning

- Recall: There is a simple way to partition a region into chunks of (approximately) equal size
- Example 20

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Partitioning By Field

- A field can be used as a coloring
- Write elements of the color space into the field **f**
 - Using an arbitrary computation
- Then call `partition(region.f, colors)`
 - Example 27

Partitioning, Digression

- Why do we want to partition data?
 - For parallelism
 - We will launch many tasks over many subregions
- A problem
 - We often need to partition multiple data structures in a consistent way
 - E.g., given that we have partitioned the nodes a particular way, that will dictate the desired partitioning of the edges

Dependent Partitioning

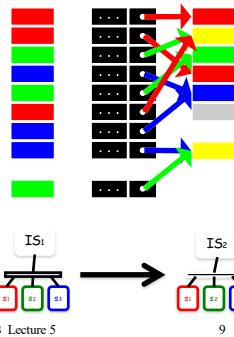
- Distinguish two kinds of partitions
- *Independent partitions*
 - Computed from the parent region, using, e.g.,
 - `partition(equals, ...)`
 - `partition(region.field, ...)`
- *Dependent partitions*
 - Computed using another partition

Dependent Partitioning Operations

- Image
 - Use the image of a field in a partition to define a new partition
- Preimage
 - Use the preimage of a field in a partition
- Set operations
 - Form new partitions using the intersection, union, and set difference of other partitions

Image

- Computes elements reachable via a field lookup
 - Equivalent to *semi-join* in relational algebra
 - Can be applied to index space or another partition
 - Computation is distributed based on location of data
- Regent understands relationship between partitions
 - Can check safety of region relation assertions at compile time

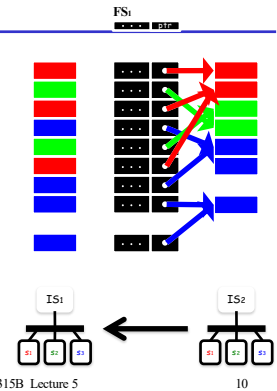


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Preimage

- Opposite of image - computes elements that reach a given subspace
 - Preserves disjointness
- Multiple images/preimages can be combined
 - can capture complex task access patterns
 - Limitation: no transitive reachability



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Example 21

- Partition the nodes
 - Equal partitioning
- Then partition the edges
 - Preimage of the source node of each edge
- For each node subregion r , form a subregion of those edges where the source node is in r

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Example 22

- Partition the edges
 - Equal partitioning
- Then partition the nodes
 - Image of the source node of each edge
- For each edge subregion r , form a subregion of those nodes that are source nodes in r

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Discussion

- Note that these two examples compute almost the same partition
- Can derive the node partition from the edges, or vice versa

Example 23

- What would the example look like if we partitioned based on the destination node?
- Let's find out ...

Set Operations: Set Difference

- Partition the edges
 - Equal partition
- Compute the source and destination node partitions of the previous two examples
- The final node partition is the set difference
 - What does this compute?
 - Examples 24 & 25

Set Operations: Set Intersection

- Partition the edges
 - Equal partition
- Compute the source & destination node partitions
- Final node partition is the intersection
 - What does this compute?
 - Example 26

Example 28

- Same as the last example
- Once the final node partition is computed, compute a partition of the edges such that each edge subregion has only the edges connecting the nodes in the corresponding node subregion

Examples 29

- Pointers point into a particular region
 - And this is part of the pointer's type
- Partitioning can change which region(s) a pointer points to
 - May lead to typechecking issues, depending on which region you want to use for an operation

Example 30

- The right way to fix type issues is to use type casts
- Very analogous to downcasting from a more general object type to a more specific object type in an object-oriented language
- But, this solution does not currently work!
 - Casting of region types not yet implemented

Example 31

- The fix/workaround is to use **wild** in field space arguments when allocating regions
- **Wild** effectively turns off typechecking for those region arguments.

Backing Up ...

- Regent's partitioning mechanisms are very different from other languages
- What do those other languages provide?

One Extreme: Simplicity

- PGAS languages (e.g. X10, UPC, Chapel) generally provide only simple array-based distribution methods
 - e.g. block, cyclic, blockcyclic
- Pros:
 - simple for programmer to describe
 - simple for compiler to verify consistency
 - simple for runtime to implement
- Cons:
 - no support for irregular (or even semi-regular) data structures
 - no support for irregular partitions of structured data
 - no support for aliased or multiple partitions

Other Extreme: Expressivity

- Initial Legion partitioning used general-purpose coloring object for ALL partitioning operations
 - Application able to color each element any way it wants
- Pros:
 - support for arbitrary irregularity in data and/or partitioning
 - support for aliased partitions, multiple partitions
- Cons:
 - significant programmer effort to describe even simple partitions
 - no ability for compiler to check that related regions are partitioned consistently
 - high runtime overhead for computing and querying partitions
 - manipulation of coloring was serial, limited to single node

Dependent Partitioning

- A carefully chosen middle ground between these two extremes
- Supports both structured and unstructured domains
- Allows arbitrary independent partitions to be computed by the application
 - But uses field data to capture intent rather than a coloring
 - Index-based partitions cover PGAS-like simple cases
- Provides an analyzable set of operations to compute dependent partitions from other partitions
 - Based on reachability and/or set operations
 - Consistency of dependent partitions can be verified at compile time
- And can be executed in parallel

Programmer Productivity

- Lines of code for computation of dependent partitions in Regent applications:

Application	Dependent		Reduction
	Original LOC	Partitioning LOC	
PENNANT	163	6	96%
Circuit	159	8	95%
MiniAero	51	7	86%

- Not a perfect metric
 - Take with however much salt you like...

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Summary

- The built-in partitioning operations are
 - Expressive
 - Can execute in parallel
 - Can be analyzed by the Regent implementation
- Except for explicit coloring objects
 - Inherently not parallel

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