

# Constructing Pairing-Friendly Elliptic Curves with Embedding Degree 10

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# Outline

- 1 Introduction
  - Pairings in Cryptography
  - Embedding Degrees: The Problem and Current Results
- 2 Constructing Curves with Prescribed Embedding Degree
  - The CM Method: The Basic Construction
  - The CM Method: Generating Families of Curves
- 3 Curves with Embedding Degree 10
  - The Construction
  - Computational Results

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# Pairings in Cryptography

- Let  $E$  be an elliptic curve defined over a finite field  $\mathbb{F}_q$ .
- For any integer  $r$  the *Weil pairing*  $e_r$  is a bilinear map sending pairs of points with order  $r$  to  $r$ -th roots of unity in  $\overline{\mathbb{F}_q}$ :

$$e_r: E[r] \times E[r] \rightarrow \mu_r.$$

- These pairings can be used in many cryptographic constructions, including:
  - one-round, 3-way key agreement (Joux);
  - short signature schemes (Boneh, Lynn, Shacham);
  - identity-based encryption (Boneh, Franklin);
  - and many others.

# Making Pairings Practical

- For pairing-based cryptosystems to be practical and secure, we require:
  - the discrete logarithm in the order- $r$  subgroup of  $E(\mathbb{F}_q)$  to be computationally infeasible;
  - the discrete logarithm in  $\mu_r$  to be computationally infeasible;
  - the pairing to be easily computable (i.e.  $\mu_r$  lies in a low-degree extension of  $\mathbb{F}_q$ ).
- To optimize applications, we want to choose a curve  $E$  so that the two discrete log problems are of about equal difficulty.

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# Embedding Degrees

- Let  $E$  be an elliptic curve over  $\mathbb{F}_q$  with  $\#E(\mathbb{F}_q) = r$ .
- Let  $k$  be the smallest integer such that  $\mu_r \subset \mathbb{F}_{q^k}^\times$   
(i.e.  $r \mid q^k - 1$ ).
  - Bal., Kob.:  $E[r] \subset E(\mathbb{F}_{q^k})$ , so the Weil pairing “embeds”  $E(\mathbb{F}_q)$  into  $\mathbb{F}_{q^k}^\times$ .
  - $k$  is the *embedding degree* of  $E$  (with respect to  $r$ ).
- Note:  $k$  is the order of  $q$  in  $(\mathbb{Z}/r\mathbb{Z})^\times$ .
  - For “random” curves,  $k \sim r \sim q$ .

# The Problem

- The problem: find primes  $q$  and elliptic curves  $E/\mathbb{F}_q$  of prime order  $r$  with small embedding degree ( $k \leq 30$ ).
  - Want to be able to control the number of bits of  $q$  to construct curves for various applications.
  - Could also look for curves  $E$  with nearly prime order –  $\#E(\mathbb{F}_q) = \text{large prime } r \times \text{small cofactor}$ .
- Requirement that  $E(\mathbb{F}_q)$  has prime order is what makes the problem difficult.
  - Cocks, Pinch: Constructed  $E$  with arbitrary embedding degree  $k$ , but largest prime factor of  $\#E(\mathbb{F}_q) \sim \sqrt{q}$ .



# Previous Results

- Menezes, Okamoto, Vanstone, 1993 (MOV): Showed that supersingular elliptic curves always have embedding degree  $k \leq 6$ .
- Miyaji, Nakabayashi, Takano, 2001 (MNT): Gave complete characterization of ordinary elliptic curves of prime order with embedding degree  $k = 3, 4, 6$ .
- Barreto, Naehrig, 2005: Constructed a family of elliptic curves of prime order and embedding degree 12.

# The Main Result

- D.F., 2006: Constructed elliptic curves of prime order with embedding degree 10.
  - Solves open problem posed by Boneh, Lynn, Shacham, 2001.

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# The CM Method of Curve Construction

- Main tool: Complex Multiplication method of curve construction (Atkin, Morain).
  - For given  $D > 0$ , constructs elliptic curve with CM by  $\mathbb{Q}(\sqrt{-D})$ .
  - Constructs curves with specified number of points.
  - Running time roughly  $O(D)$ .
- How it works: Fix  $D, k$ . Look for  $t, n, q$  (representing trace, number of points, and size of field) satisfying
  - 1  $q, n$  prime;
  - 2  $n = q + 1 - t$  (formula for number of points);
  - 3  $n$  divides  $q^k - 1$  (embedding degree  $k$ );
  - 4  $Dy^2 = 4q - t^2$  for some integer  $y$ .
- For such  $t, n, q$ , if  $D$  is not too large ( $\sim 10^9$ ) we can construct an elliptic curve  $E$  over  $\mathbb{F}_q$  with  $n$  points and embedding degree  $k$ .

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# Observations about the CM Method

- Barreto, Lynn, Scott: The embedding degree condition  $n \mid q^k - 1$  can be replaced with  $n \mid \Phi_k(t - 1)$ , where  $\Phi_k$  is the  $k$ -th cyclotomic polynomial.
  - $k$  smallest such that  $n \mid q^k - 1$  implies  $n \mid \Phi_k(q)$ .
  - $q + 1 - t = n$  implies  $q = t - 1 \pmod{n}$ .
- Barreto, Naehrig: Parametrize  $t, n, q$  as polynomials:  $t(x), n(x), q(x)$ . Construct curves by finding integer solutions to

$$Dy^2 = 4q(x) - t(x)^2 = 4n(x) - (t(x) - 2)^2.$$

# Extending the CM Method

- Combine these observations to get an algorithm for generating families of pairing-friendly curves:
  - 1 Fix  $D$ ,  $k$ , and choose a polynomial  $t(x)$ .
  - 2 Choose  $n(x)$  an irreducible factor of  $\Phi_k(t(x) - 1)$ .
  - 3 Find integer solutions  $(x, y)$  to  $Dy^2 = 4n(x) - (t(x) - 2)^2$ .
  - 4 If  $q(x)$ ,  $n(x)$  are both prime, construct elliptic curve over  $\mathbb{F}_{q(x)}$  with  $n(x)$  points.
- Step 3 is the difficult part: if  $f(x) = 4n(x) - (t(x) - 2)^2$  has degree  $\geq 3$ , then  $Dy^2 = f(x)$  (usually) has only a finite number of integer solutions! (Siegel's Theorem)

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# Embedding Degree 10: The Key Observation

- Goal: Choose  $t(x)$ , choose  $n(x)$  an irreducible factor of  $\Phi_{10}(t(x) - 1)$ , such that  $4n(x) - (t(x) - 2)^2$  is quadratic.
  - $\Phi_{10}(x) = x^4 - x^3 + x^2 - x + 1$  degree 4.
  - All irred. factors of  $\Phi_{10}(t(x) - 1)$  must have  $4 \mid \text{degree}$ .
- Key observation: Need to choose  $n(x)$ ,  $t(x)$  such that the leading terms of  $4n$  and  $t^2$  cancel out.
  - Smallest possible case:  $\text{deg } n = 4$ ,  $\text{deg } t = 2$ .
- Galbraith, McKee, Valença: Characterized quadratic  $t(x)$  such that  $\Phi_{10}(t(x) - 1)$  factors into two quartics.
- One of these  $t(x)$  gives the desired cancellation!

# Choice of Parameters

- Choose  $t, n, q$  as follows:

$$t(x) = 10x^2 + 5x + 3$$

$$n(x) = 25x^4 + 25x^3 + 15x^2 + 5x + 1$$

$$q(x) = 25x^4 + 25x^3 + 25x^2 + 10x + 3$$

- Then  $n(x)$  divides  $\Phi_{10}(t(x) - 1)$ , and

$$f(x) = 4n(x) - (t(x) - 2)^2 = 15x^2 + 10x + 3.$$

# Constructing the Curves

- To find curves: set  $D' = 15D$ ,  $u = 15x + 5$ , complete the square in  $Dy^2 = 15x^2 + 10x + 3$  to get

$$u^2 - D'y^2 = -20.$$

- Find integer solutions  $(u, y)$  to this Pell-like equation. (Use LMM algorithm.)
- Let  $x = (u - 5)/15$ . If  $q(x)$  and  $n(x)$  are prime, there exists an elliptic curve with embedding degree 10!
- If no solutions of appropriate size, or  $q(x)$  or  $n(x)$  not prime, increase  $D$  and try again.

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# General Results (Thanks to Mike Scott)

- Searched for curves with  $D < 2 \cdot 10^9$ ,  $q$  between 148 and 512 bits.
  - Crypto applications: want  $q \sim 220$ -250 bits (discrete log in  $E(\mathbb{F}_q)$  about as difficult as discrete log in  $\mathbb{F}_{q^{10}}$ ).
- Found 23 curves of prime order with embedding degree 10.
- Found 101 curves of nearly prime order (large prime  $\times$  small cofactor  $< 2^{16}$ ) and embedding degree 10.
- Ability to handle larger  $D$  in CM Method would allow us to find more curves.

# Example: A 234-bit Curve (Computed by Mike Scott)

- Set  $D = 1227652867$ .
- Compute solution  $(x, y)$  to  $Dy^2 = 15x^2 + 10x + 3$ .
- Use this value of  $x$  to compute

$$t = 269901098952705059670276196260897153$$

$$n = 18211650803969472064493264347375949776033155743952030750450033782306651$$





$$q = 18211650803969472064493264347375950045934254696657090420726230043203803$$

- Use CM method to compute curve equation over  $\mathbb{F}_q$ :  
(Given  $t, n, q$ , curve equation took about a week to compute.)




$$y^2 = x^3 - 3x + 15748668094913401184777964473522859086900831274922948973320684995903275.$$

- This curve has  $n$  points and embedding degree 10!

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