PARTIAL AND COMPLETE CYCLIC ORDERS

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We show that, in contrast to a famous theorem on linear orders, not every partial cyclic order on $M = \{1, \ldots, m\}$ can be extended to a complete cyclic order. In fact, the complexity, in a certain sense, of sufficient conditions for such an extendability increases rapidly with m.

DEFINITION 1. (i) Two linear orders, (a_1, \ldots, a_m) and (b_1, \ldots, b_m) , on M are called *cyclically equivalent* if there exists $k \in M$ such that $[j-1 \equiv (i-1+k) \pmod{m}] \Rightarrow a_i = b_i$.

(ii) A complete cyclic order (CCO) on M is an equivalence class C of linear orders modulo cyclic equivalence; denote $a_1 a_2 \cdots a_m$ for the equivalence class containing (a_1, a_2, \ldots, a_m) .

DEFINITION 2. A partial cyclic order (PCO) on M is a set Δ of cyclically ordered triples (COTs) out of M such that:

- (i) $xyz \in \Delta \Rightarrow zyx \notin \Delta$ ("antisymmetry"),
- (ii) $\{xyz, xzw\} \subset \Delta \Rightarrow xyw \in \Delta$ ("transitivity"); since xyz = zxy, etc., also $yzw \notin \Delta$ is implied.

THEOREM 3. (i) If C is a CCO then the set Δ of all COTs derived from C is a PCO. (ii) If Δ is a saturated PCO, i.e., $\{x, y, z\} \in \binom{M}{3}$ & $xyz \notin \Delta \Rightarrow zyx \in \Delta$, then there exists a CCO from which all of Δ 's COTs are derived; Δ is then said to be extendable to a CCO.

COROLLARY 4. A PCO is extendable to a CCO if and only if it is contained in a saturated PCO.

It is natural to ask whether every PCO is extendable to a CCO (or, equivalently, is contained in a saturated PCO). In view of the following example, the answer is in the negative.

EXAMPLE 5. Let $M = \{a, b, \ldots, m\}$ be the set of the first thirteen letters, and let $\Delta = \{acd, bde, cef, dfg, egh, fha, gac, hcb, abi, cij, bjk, ikl, jlm, kma, lab, mbc, hcm, bhm}. Obviously, <math>\Delta$ is a PCO. Suppose that $\Delta^* \supset \Delta$ is a saturated PCO. If $abc \in \Delta^*$ then, since $acd \in \Delta^*$, also $bcd \in \Delta^*$. Then, also $cde \in \Delta^*$, and successive applications of transivity finally yield $acb \in \Delta^*$, which contra-

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dicts antisymmetry. Thus, $abc \notin \Delta^*$ and, therefore, by saturatedness, $acb \in \Delta^*$. Analogously, since $abi \in \Delta^{*_1}$, also $cbi \in \Delta^*$, and successive applications of transivity finally yield $abc \in \Delta^*$. Thus, antisymmetry is contradicted again. It follows that there is no saturated PCO that contains Δ .

The unextendability of Δ in Example 5 followed essentially from the fact that neither *abc* nor *cba* belongs to any PCO that contains Δ . That gives rise to the following definition.

DEFINITION 6. If $\tau = \{i, j, k\}, 1 \le i < j < k \le m$, denote $\tau^+ = ijk$ and $\tau^- = kji$ for the two possible cyclic orderings of τ . A PCO Δ is said to satisfy the nth order conditon if for every $\tau_1, \ldots, \tau_n \in \binom{M}{3}$ there exists a PCO $\Delta^* \supset \Delta$ and $\epsilon_i \in \{+, -\}$ $(i = 1, \ldots, n)$, such that $\{\tau_i^{\epsilon_1}, \ldots, \tau_n^{\epsilon_n}\} \subset \Delta^*$.

Obviously, all the *n*th order conditions (n = 0, 1, ...) are necessary for extendability to a CCO and, as *n* increases, the *n*th order condition becomes stronger. The conjunction of all the *n*th order conditions (n = 0, 1, ...) is a sufficient condition for every *m*. It is natural to ask whether there exists an *n* such that the *n*th order condition suffices for every PCO Δ on a finite set *M* to be extendable at a CCO. Unfortunately, the answer to this question also is in the negative. A sequence of PCOs that prove this is constructed as follows.

EXAMPLE 7. Let $m_0=13$ and let Δ_0 be the PCO on $M_0=\{1,\ldots,13\}$ defined in Example 5 (identify a with 1, b with 2, etc.). As we have already seen, Δ_0 is not extendable to a CCO. However, since it is a PCO, it satisfies the 0th order condition. For the purpose of later use in induction, note that $\Delta_0\setminus\{egh\}$ is extendable to the following complete cyclic ordering: afbhcgdeijklm. Suppose, by induction, that Δ_n is a PCO on $M_n=\{1,\ldots,m_n\}, \Delta_n$ satisfies the the nth order condition but is not extendable to a CCO. Suppose also that $xyz\in\Delta_n$ is such that $\Delta_n\setminus\{xyz\}$ is extendable to a CCO. We construct Δ_{n+1} as follows. Define $m_{n+1}=m_n+15$ and $M_{n+1}=\{1,\ldots,m_{n+1}\}$ and denote $(u_1,\ldots,u_5,v_1,\ldots,v_5,w_1,\ldots,w_5)=(m_n+1,\ldots,m_{n+1})$. Let

$$\begin{split} \Delta' &= (\Delta_n \backslash \{xyz\}) \ \cup \ \{zu_1u_2, yu_2u_3, u_1u_3u_4, u_2u_4u_5, u_3u_5x, u_4xy, xyv_1, \\ & u_1v_1v_2, yv_2v_3, v_1v_3v_4, v_2v_4v_5, v_3v_5x, v_4xy, v_5yu_1, \\ & u_5yw_1, zw_1w_2, yw_2w_3, w_1w_3w_4, w_2w_4w_5, w_3w_5u_5, \\ & w_4u_5y, w_5y^2\}. \end{split}$$

Let Δ_{n+1} be the transitive closure of Δ' , i.e., Δ_{n+1} is the intersection of all the transitive classes of COTs that contain Δ' (see Definition 2). It turns out that Δ_{n+1} is a PCO, but is not extendable to a CCO. Also, $\Delta_{n+1} \setminus \{w_3w_5u_5\}$ is extendible to a CCO. The proof of these facts follows from the analogous properties of Δ_n . The important property of Δ_{n+1} is that it satisfies the (n+1)st order condition. A detailed proof will be given elsewhere. Here, we indicate that two cases are distinguished when a set of 3-element subsets of M_{n+1} is given.

First, when $|\tau_i \cap M_{n+1}| \ge 2$ for $i=1,\ldots,n+1$, the ϵ_i-s are determined essentially by the CCO on M_n to which $\Delta_n \setminus \{xyz\}$ is extendible. Otherwise, the induction hypothesis is applied and the ϵ_i-s are determined essentially by a PCO Δ^* that contains Δ_n and n of the τ_i-s .

In view of Example 7, an algorithm for extending a PCO to a CCO which is based on successive addings of COTs, cannot be polynomial. We conjecture that there is no polynomial algorithm for this problem; note that there seems to be an equivalence between our problem and that of the Hamiltonian path, from the point of view of complexity of computations.

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