

## Problem Set 3

1. Let  $G : \{0, 1\}^n \rightarrow \{0, 1\}^{2n}$  be a  $(t, \epsilon)$ -secure pseudorandom generator computable in time  $r$ .

Show that  $G$  is also a  $(t - r - O(n), \epsilon + 2^{-n})$ -secure one way function.

2. Let  $f : \{0, 1\}^n \rightarrow \{0, 1\}^m$  be a  $(t, \epsilon)$ -secure one-way function.

Show that

$$\frac{t}{\epsilon} \leq O((m + n) \cdot 2^n)$$

3. Let  $f : \{0, 1\}^n \rightarrow \{0, 1\}^n$  be a  $(t, \epsilon)$ -secure one-way permutation computable in time  $\leq r$ .

Show that

$$\frac{t^2}{\epsilon} \leq O((r + n^2)^2 \cdot 2^n)$$

[Hint: first show that, for any permutation  $f : \{0, 1\}^n \rightarrow \{0, 1\}^n$ , there is an algorithm of complexity  $O(r \cdot 2^{n/2})$  that inverts the permutation everywhere. The algorithm is given a pre-computed data structure of size  $O(n2^{n/2})$  and runs in time  $O(r2^{n/2})$ . Recall that in our model of computation we do not pay for the price of pre-computing data at “compile time,” we only pay the sum of the length of the program, including any fixed data it needs access to, plus the worst-case running time.]

4. Let  $f : \{0, 1\}^n \rightarrow \{0, 1\}^n$  be a  $(t, \epsilon)$ -secure one-way function computable in time  $r$ .

Show that  $g : \{0, 1\}^{2n} \rightarrow \{0, 1\}^{2n}$  defined as  $g(x, y) := f(x), f(y)$  is  $(t - O(r), \epsilon)$ -secure.

5. Let  $p$  be a prime and  $g$  be a generator for  $\mathbb{Z}_p^*$  such that  $f(x) := g^x \bmod p$  is a  $(t, 0.99)$ -one way permutation. Let  $k = \lceil \log_2 p \rceil$  be the number of digits of  $p$ ; then recall that  $f()$  is computable in time  $O(k^3)$ .

Show that  $f()$  is also  $(\frac{1}{24,000}(t - O(k^3)), 0.51)$ -one way.

6. Recall that if  $F : \{0, 1\}^n \rightarrow \{0, 1\}^n$  is a function then we define the Feistel permutation  $D_F(x, y) := y, x \oplus F(y)$ .

Show that there is an efficient oracle algorithm  $A$  such that

$$\mathbb{P}_{\Pi: \{0,1\}^{2m} \rightarrow \{0,1\}^{2m}} [A^{\Pi, \Pi^{-1}} = 1] = 2^{-\Omega(m)}$$

where  $\Pi$  is a random permutation, but for every three functions  $F_1, F_2, F_3$ , if we define  $P(x) := D_{F_3}(D_{F_2}(D_{F_1}(x)))$  we have

$$A^{P, P^{-1}} = 1$$

[Note: I don't know the solution.]