

# CS265/CME309: Randomized Algorithms and Probabilistic Analysis

## Lecture #16: Martingales, the Doob Martingale, and Azuma-Hoeffding Tail Bounds

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### 1 Introduction

In this lecture we introduce the concept of a Martingale, which can be used to analyze many random processes. The technical core of this lecture will be the Azuma-Hoeffding tail bound, which gives a Chernoff style bound on tail probabilities of martingales. These bounds were independently proved by Azuma and Hoeffding in the 1960's [1, 2]. Crucially, these bounds will let us often effortlessly prove strong concentration results for sums of certain *dependent* random variables and other stochastic processes. The proof technique for the Azuma-Hoeffding bounds follow the same general approach as Chernoff bounds: applying Markov's inequality to the moment generating function for the random variable in question. Our analysis of this moment generating function, however, will crucially leverage the properties of martingales, as opposed to the properties of sums of independent random variables.

**Definition 1.** *A sequence of real-valued random variables,  $\{Z_t\}$  is a martingale with respect to the sequence (of not necessarily real valued random variables)  $\{X_t\}$  if the following conditions hold for all  $t$ :*

1.  $Z_t$  is a function of  $X_1, \dots, X_t$ ,
2.  $\mathbf{E}[|Z_t|] < \infty$ ,
3.  $\mathbf{E}[Z_t | X_0, \dots, X_{t-1}] = Z_{t-1}$ .

Given this setting, letting  $Y_t = Z_t - Z_{t-1}$ , often refers to the sequence  $\{Y_t\}$  as the sequence of martingale differences. In many cases,  $\{X_t\}$  is a martingale with respect to itself, in the sense that  $\mathbf{E}[X_t | X_0, \dots, X_{t-1}] = X_{t-1}$ . In such cases, one generally just says that  $\{X_t\}$  “is a martingale”.

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The above definition can be slightly generalized via the notion of *filters* and  $\sigma$ -fields, though we will not get bogged down in that. For our purposes, the above definition will be sufficient. The first condition, that  $Z_t$  is a function of  $X_1, \dots, X_t$  is not too restrictive: if we would like  $Z_t$  to be a randomized function, we can regard  $Z_t$  as being a deterministic function of  $X_1, \dots, X_t$  in addition to some random bits  $b_0^t, \dots, b_m^t$  and then define  $Z_t$  as a martingale with respect to the sequence  $\{X'_t\}$  where  $X'_t = (X_t, b_0^t, \dots, b_m^t)$ . Hence we can assume that  $Z_t$  is a deterministic function of the  $X_0, \dots, X_t$ . The above definition can be applied both when we have a finite set of  $Z_t$ 's, as well as the infinite case.

## 1.1 The Doob Martingale

One reason why martingales are such a useful concept is that we can obtain a martingale from *any* real-valued random variable.

**Definition 2.** *Given a random variable,  $A$ , and a sequence of random variables,  $\{X_t\}$  defined over a common probability space (i.e. quantities such as  $\Pr[A = 1 | X_4 = 7]$  are well defined notions), the Doob martingale of  $A$  with respect to  $\{X_t\}$  is the sequence  $\{Z_t\}$  defined by*

$$Z_t = \mathbf{E}[A | X_0, \dots, X_t].$$

To see why the above definition yields that  $\{Z_t\}$  is a martingale with respect to  $\{X_t\}$ , provided  $\mathbf{E}[|Z_t|] < \infty$ , we simply apply the definition of  $Z_t$  and simplify:

$$\mathbf{E}[Z_t | X_0, \dots, X_{t-1}] = \mathbf{E}[\mathbf{E}[A | X_0, \dots, X_t] | X_0, \dots, X_{t-1}] = \mathbf{E}[A | X_0, \dots, X_{t-1}] = Z_{t-1}.$$

In the above, we leveraged the fact that for any random variables  $X, Y, Z$ , by the definition of conditional probability,

$$\mathbf{E}[\mathbf{E}[X | Y, Z] | Y] = \mathbf{E}[X | Y],$$

since the conditioning on  $Z$  is integrated/marginalized out by the outer expectation.

One of the most common ways that the Doob martingale is employed is to analyze properties of some random variable that depends on a random object—for example, the number of empty bins if we randomly toss  $m$  balls into  $n$  bins, or the chromatic number of an Eros-Renyi random graph, etc. In those settings, the  $X_t$ 's sequentially reveal information about the random object and the quantity in question, until there is no more randomness. The following examples illustrate this use-case of the Doob martingale.

**Example 3.** *Consider tossing  $m$  balls uniformly at random into  $n$  bins. Let  $A$  denote the number of empty bins at the end of the process. Letting  $X_t$  denote the location that the  $t$ th ball lands, consider the Doob martingale defined by  $Z_t = \mathbf{E}[A | X_1, \dots, X_t]$ . Here, we are iteratively updating our expectation of  $A$  as we see where each of the  $n$  balls is thrown.  $Z_0 = \mathbf{E}[A]$ , and  $Z_n$  equals the actual value of  $A$  for the realization of the ball-tossing process given by the  $X_t$ 's.*

**Example 4.** *Consider the Erdos-Renyi random graph process  $G_{n,p}$  that constructs an  $n$  node graph by independently including each edge with probability  $p$ . Let  $A$  denote the chromatic number of the resulting graph, defined as the smallest number of colors such that it is possible to color the nodes in such a way that no neighboring nodes have the same color. If we define the random variables  $X_1, \dots, X_{\binom{n}{2}}$  such that  $X_t$  is the indicator for whether the  $t$ th edge is present in the graph, then we*

have the corresponding Doob martingale defined by  $Z_t = \mathbf{E}[A|X_1, \dots, X_t]$ . This is known as the edge exposure martingale, as we are revealing the edges one-at-a-time.

Alternately, we could define the random variables  $Y_1, \dots, Y_n$ , such that  $Y_t$  is the random variable describing the neighborhood of the  $t$ th node, namely this random variable takes one of  $2^{n-1}$  different values, according to which of the  $n-1$  potential edges are present between node  $t$  and the  $n-1$  other nodes. We can now define the vertex exposure martingale  $Z_1, \dots, Z_n$  with  $Z_t = \mathbf{E}[A|Y_1, \dots, Y_t]$ .

## 2 The Azuma-Hoeffding Tail Bound

**Theorem 1.** Let  $\{Z_t\}$  be a martingale with respect to  $\{X_t\}$ , and suppose there are constants  $c_1, \dots, c_n$  such that for all  $i \leq n$ ,  $|Z_i - Z_{i-1}| \leq c_i$ . For any  $\lambda > 0$ ,

$$\Pr[|Z_n - Z_0| \geq \lambda] \leq 2e^{-\frac{\lambda^2}{2\sum c_i^2}}.$$

To compare this with the standard Chernoff bounds, note that if  $X_1, \dots, X_n$  denote the 0/1 outcomes of  $n$  coin flips, and  $Z_t = \mathbf{E}[\sum X_i | X_1, \dots, X_t]$ , is the Doob martingale corresponding to the total number of heads out of the  $n$  tosses, then the above theorem applies with  $c_i = 1/2$ , and we get  $\Pr[|Z_t - n/2| \geq c\sqrt{n}] \leq 2e^{-\frac{c^2 n}{2n \cdot (1/4)}} = 2e^{-2c^2}$ , which is our familiar inverse exponential tail bounds.

The following lemma will be central to our proof of Theorem 1.

**Lemma 5.** Let  $Y$  denote a random variable  $Y \in [-c, c]$  for some constant  $c$ , with  $\mathbf{E}[Y] = 0$ , then for all  $t \geq 0$ ,

$$\mathbf{E}[e^{tY}] \leq e^{t^2 c^2 / 2}.$$

*Proof.* We will first prove the statement when  $c = 1$ , from which the claim will follow by observing that if  $Y \in [-c, c]$ , the random variable  $Y' = Y/c$  is in the interval  $[-1, 1]$ , and  $tY = (tc)Y'$ , hence  $\mathbf{E}[e^{tY}] = \mathbf{E}[e^{(tc)Y'}]$ . So if we show that  $\mathbf{E}[e^{tY'}] \leq e^{t^2/2}$ , the claim will follow.

We now analyze  $\mathbf{E}[e^{tY}]$  for a random variable  $Y$  with expectation 0 and magnitude bounded by 1. From the convexity of the exponential function, for any  $x \in [-1, 1]$ ,

$$e^{tx} \leq \frac{(1-x)e^{t(-1)} + (1+x)e^{t(1)}}{2}.$$

Hence

$$\mathbf{E}[e^{tY}] \leq \mathbf{E}\left[\frac{(1-Y)e^{-t} + (1+Y)e^t}{2}\right] = \frac{e^{-t} + e^t}{2},$$

where the final equality is from the assumption that  $\mathbf{E}[Y] = 0$ . To finish the proof, we will simply re-express  $e^{-t}$  and  $e^t$  in terms of their Taylor expansions, and compare this to the Taylor expansion of  $e^{t^2/2}$ . To this end, we have

$$\frac{1}{2}(e^{-t} + e^t) = \frac{1}{2} \left( 1 + t + \frac{t^2}{2} + \frac{t^3}{3!} + \dots + 1 - t + \frac{t^2}{2} - \frac{t^3}{3!} + \dots \right) = 1 + \frac{t^2}{2} + \frac{t^4}{4!} + \frac{t^6}{6!} + \dots$$

On the other hand, the Taylor expansion of  $e^{t^2/2}$  is the following:

$$e^{t^2/2} = \sum_{j \geq 0} \frac{x^2 j}{2^j j!} = 1 + \frac{t^2}{2} + \frac{t^4}{4 \cdot 2} + \frac{t^6}{2^3 \cdot 3!} + \dots,$$

hence  $\frac{1}{2}(e^{-t} + e^t) \leq e^{t^2/2}$  because each term of the expansion on the left is at least the corresponding term in the expansion of the right (and both sequences converge).  $\square$

We now prove Theorem 1.

*Proof of Theorem 1.* We begin with the upper bound. For any  $t > 0$ , by the monotonicity of the exponential, and Markov's inequality,

$$\Pr[Z_n - Z_0 \geq \lambda] = \Pr[e^{t(Z_n - Z_0)} \geq e^{t\lambda}] \leq \frac{\mathbf{E}[e^{t(Z_n - Z_0)}]}{e^{-t\lambda}}.$$

Note that for any random variables  $A, B$ ,  $\mathbf{E}[A] = \mathbf{E}[\mathbf{E}[A|B]]$ , hence

$$\mathbf{E}[e^{t(Z_n - Z_0)}] = \mathbf{E}[e^{t(Z_n - Z_{n-1} + Z_{n-1} - Z_0)}] = \mathbf{E}[\mathbf{E}[e^{t(Z_n - Z_{n-1} + Z_{n-1} - Z_0)} | X_0, \dots, X_{n-1}]] .$$

Since  $Z_{n-1}$  is a function of  $X_0, \dots, X_{n-1}$  we can pull the “constant”  $e^{t(Z_{n-1} - Z_0)}$  out from the inner expectation, yielding:

$$\mathbf{E}[\mathbf{E}[e^{t(Z_n - Z_0)} | X_0, \dots, X_{n-1}]] = \mathbf{E}[e^{t(Z_{n-1} - Z_0)} \mathbf{E}[e^{t(Z_n - Z_{n-1})} | X_0, \dots, X_{n-1}]] \leq \mathbf{E}[e^{t(Z_{n-1} - Z_0)}] e^{t^2 c_n^2 / 2},$$

where the last inequality resulted from applying Lemma 5 to the random variable  $Z_n - Z_{n-1} | X_0, \dots, X_{n-1}$  which is zero in expectation and has magnitude bounded by  $c_n$ .

Iteratively applying the above process, we get that

$$\mathbf{E}[e^{t(Z_n - Z_0)}] \leq e^{\frac{t^2 \sum_{i=1}^n c_i^2}{2}} .$$

Hence putting the pieces together, and plugging in  $t = \frac{\lambda}{\sum c_i^2}$ , we obtain the following:

$$\Pr[Z_n - Z_0 \geq \lambda] \leq \frac{\mathbf{E}[e^{t(Z_n - Z_0)}]}{e^{-t\lambda}} \leq e^{\frac{t^2 \sum_{i=1}^n c_i^2}{2}} e^{-t\lambda} = e^{-\frac{\lambda^2}{2 \sum c_i^2}} .$$

The lower tail bound is proved analogously.  $\square$

## 2.1 Applications of Azuma-Hoeffding

In many settings, the Azuma-Hoeffding tail bound will be easy to apply, because the underlying quantity we care about corresponds to a clean martingale. It is, however, always worth stopping and thinking about whether there is a *better* martingale that you could use—better in the sense that the bounds on the martingale differences,  $c_i$ , are smaller, or better in the sense that there are fewer steps in the sequence (e.g. the difference between the  $\binom{n}{2}$  length edge-exposure martingale of Example 4 versus the  $n$  terms in the vertex-exposure martingale).

**Proposition 6.** *Letting  $A$  denote the chromatic number of an Erdos-Renyi random graph  $G_{n,p}$ ,*

$$\Pr[|A - \mathbf{E}[A]| \geq c\sqrt{n}] \leq e^{-\frac{c^2}{2}} .$$

*Proof.* Recalling the vertex exposure martingale from Example 4, it is clear that  $Z_t - Z_{t-1} \leq 1$ , since revealing the neighborhood of a vertex can change the expected chromatic number by at most 1 (in the best case, we save using a color, in the worst case, we need a new color just for that vertex). Hence Theorem 1 holds with  $c_i = 1$ , and the proposition holds. [Note that if we had used the edge exposure martingale instead, we would have gotten the much worse bound of  $\Pr[|A - \mathbf{E}[A]| \geq c\sqrt{n}] \leq e^{-\frac{c^2}{n}}$ , since there we still would have had  $c_i = 1$ , but we would have had  $\binom{n}{2} \approx n^2/2$  terms in our sequence.  $\square$ ]

The next example illustrates how one can deal with martingales that depend on additional randomness beyond the natural choice of  $\{X_t\}$  by defining an alternate sequence  $\{X'_t\}$  where each  $X'_t$  includes both  $X_t$  as well as some additional randomness.

**Example 7.** Consider a gambling setting where, at each timestep, you can bet any amount between 0 and some “house limit”  $B$ . At each step, a fair coin is flipped, and if you correctly guess the outcome, you win your bet, otherwise you lose your bet (so the expected earnings in any round are 0). Let  $\{X_t\}$  denote the outcome of the coin flip in each round. Intuitively, since the expected winnings at each step is 0, letting  $\{Z_t\}$  denote our sequence of net earnings (i.e.  $Z_t$  is our net winnings or losses after  $t$  rounds), because  $|Z_t - Z_{t-1}| \leq B$ , we can apply the Azuma-Hoeffding bound to conclude that

$$\Pr[|Z_t| \geq \lambda] \leq 2e^{-\frac{\lambda^2}{2tB^2}}.$$

Crucially, the above bound holds no matter what betting strategy we use. Even if our bet in the  $t$ th round depends on the outcomes in the previous  $t - 1$  rounds, this still holds. [This is why casino owners can sleep soundly at night—they set the house limit  $B$  to be low enough, in comparison to the number of rounds that will be played, and the small but non-zero bias that is build into the games, so that the probability that they don’t end up making money in a given week is extremely small. The inverse exponential with  $\lambda^2$  in the exponent makes it okay for them to use a fairly large value  $B$ . . . .]

One final note: if we want to be super pedantic and make sure that  $\{Z_t\}$  is actually a martingale, even for a randomized betting strategy that depends on previous rounds, and on other random whims, etc., then we would say that  $\{Z_t\}$  is a martingale with respect to the sequence  $\{X'_t\}$ , where  $X'_t$  denotes both the outcome of the  $t$ th round, as well as the bet we will place in the  $t + 1$ st round. This is a perfectly valid sequence of random variables, and it is true that  $\{Z_t\}$  is a martingale with respect to  $\{X'_t\}$ , since  $Z_t$  is a deterministic function of  $X'_0, \dots, X'_t$  and  $\mathbf{E}[Z_{t+1}|X'_0, \dots, X'_t] = Z_t$ .]

## References

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- [2] Wassily Hoeffding. Probability inequalities for sums of bounded random variables. *Journal of the American Statistical Association*, 58(301):13–30, 1963.