

Conclusions and Future Directions

Amir Abboud

Stanford University



“Hardness in P” workshop, STOC 15’

This talk:

➔ Quick recap.

➔ What's next for Hardness in P?

Five of the many possible directions I am excited about.

Take a problem X in P , say in $O(n^2)$ time.

And prove that:

“ X probably *cannot* be solved in $O(n^{2-e})$ time.”

We imitate NP-hardness.

An $O(n^{1.9})$ algorithm
for problem X

reductions
→

A popular and plausible
conjecture is falsified

Three popular conjectures

The 3-SUM Conjecture:

“No $O(n^{2-\epsilon})$ time algorithm for finding three numbers that sum to 0 in a list of n integers.”

The APSP Conjecture:

“No $O(n^{3-\epsilon})$ time algorithm for computing all pairs shortest paths in an edge weighted graph on n nodes.”

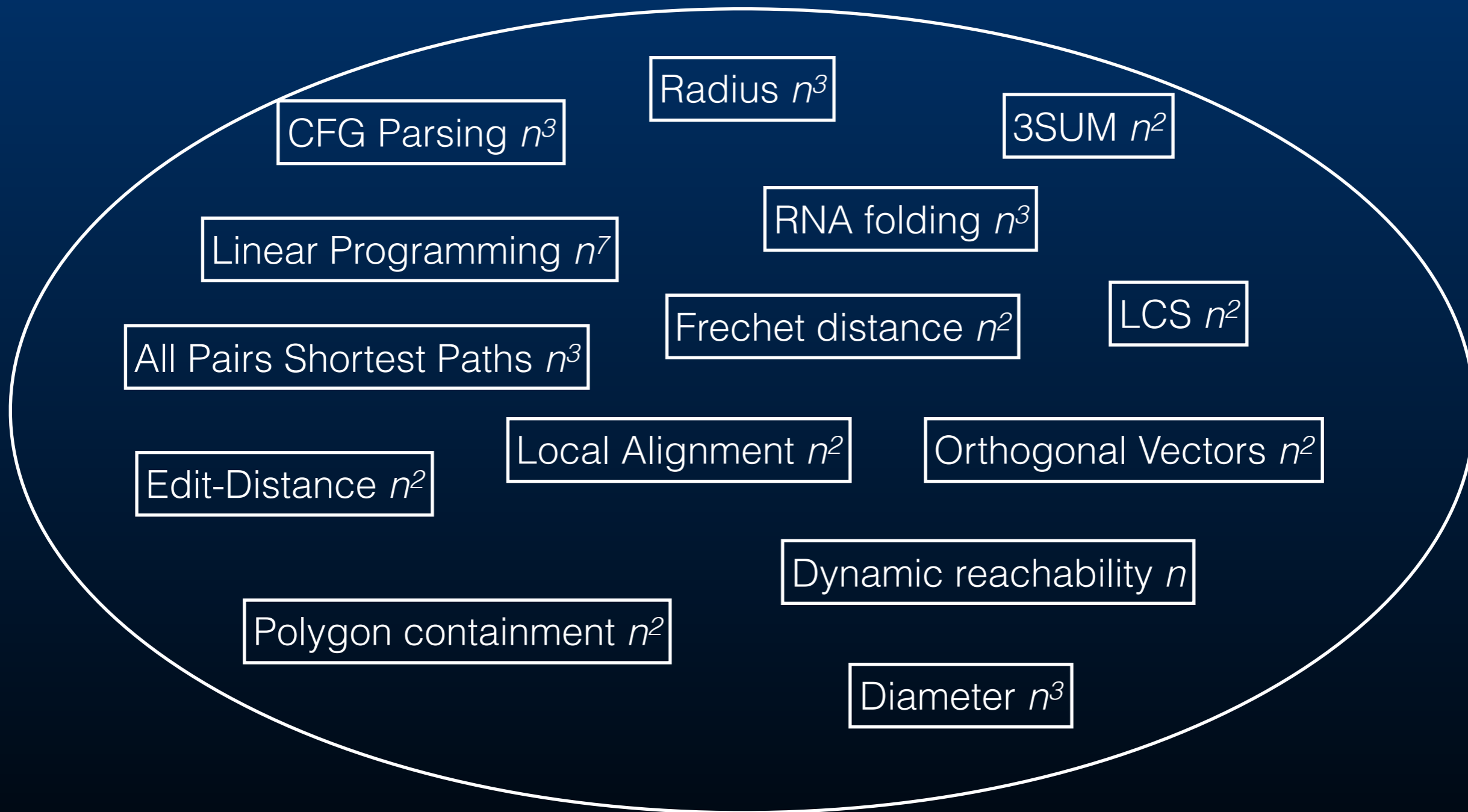
The Orthogonal Vectors Conjecture:

(implied by the Strong Exponential Time Hypothesis)

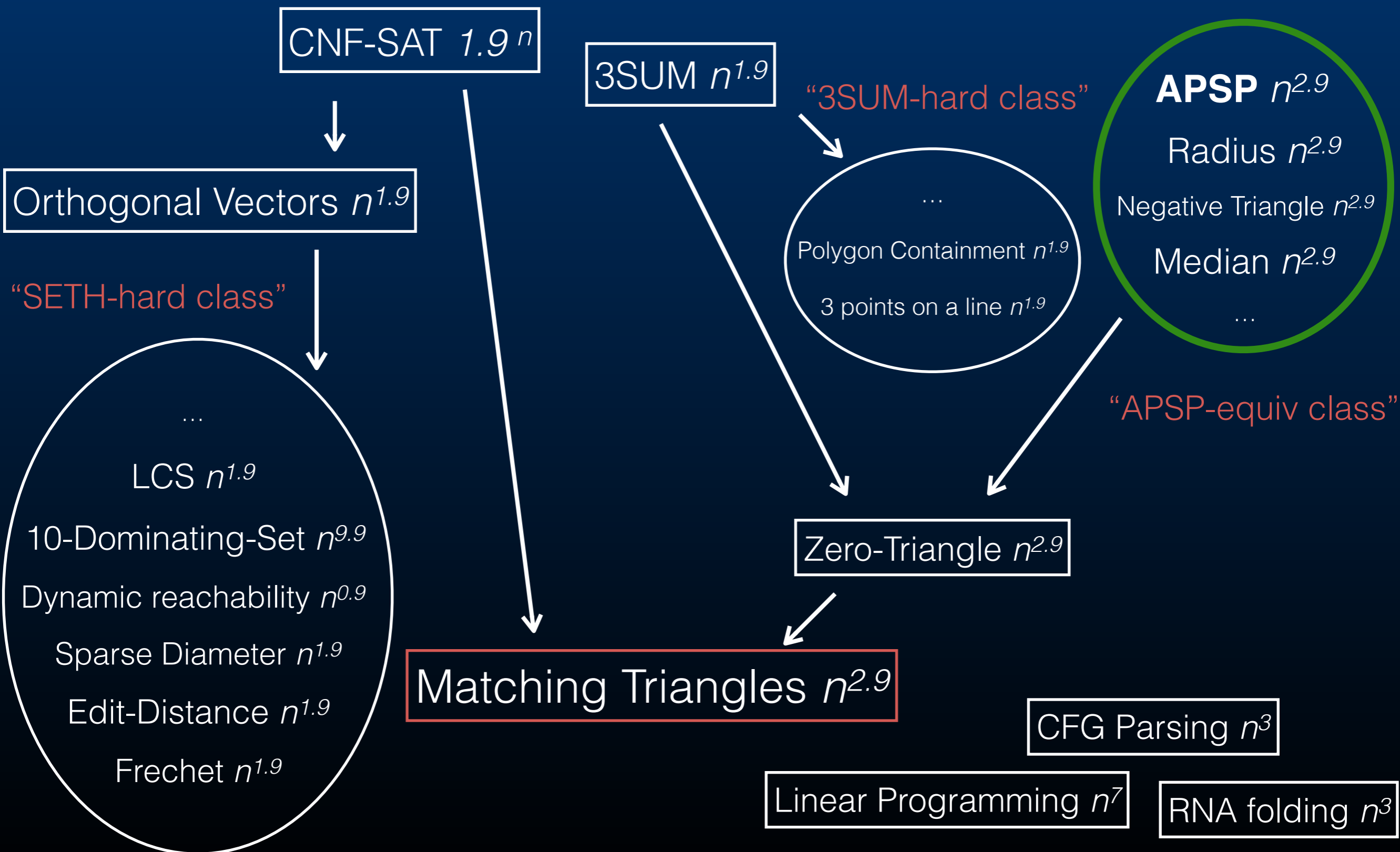
“No $O(n^{2-\epsilon})$ time algorithm for finding a pair of orthogonal vectors in a list of n boolean vectors of length $\log^2 n$.”

Tight lower bounds for diverse problems!

P before...



P after...



Many intriguing questions remain!

➔ What's next for "Hardness in P"?

- Better understanding of the conjectures

"strengthening the foundations of this project"

Find more reasons to believe the conjectures

“If 3SUM is in $n^{1.9}$ time, then...”

“If CNF-SAT is in 1.9^n time, then...”

Alternatively, replacing the conjectures
with more plausible ones.

[A-VW-Yu STOC 15'] tomorrow!

Find more reasons to **disbelieve** the conjectures?

Very interesting progress on 3SUM.

[Chan-Lewenstein STOC 15'] tomorrow!

Find connections between the conjectures

“ maybe all $O(n^2)$ problems are equivalent under subquadratic reductions ”



big open questions.

Find barriers for relating them?

➔ What's next for "Hardness in P"?

- Better understanding of the conjectures
 - Classifying more problems

"If we're stuck with a bound, we should know why."

To classify all of P, new conjectures might be needed.

Unclassified problems:

Maximum Matching

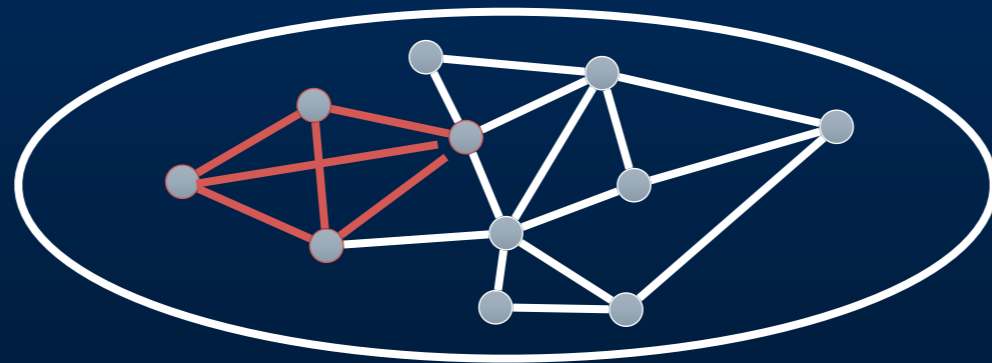
Linear Programming

Can we write a small LP for Orthogonal Vectors or 3SUM?

Recently: a new class based on k-Clique.

k-Clique

Given a graph on n nodes, are there k that form a clique?



Best combinatorial algorithm: $\sim O(n^k)$

[Nesetril - Poljak 85']: $O(n^{\omega k/3}) = O(n^{0.79k})$

[A-Backurs-VW 15']

If these algorithms are optimal, we can show very nice lower bounds in P!

[A-Backurs-VW Arxiv 15'] k-Clique based lower bounds

A very basic CS problem

CFG Parsing

Input: Context Free Grammar G and a string w of length n

G

$S \rightarrow AB$
$A \rightarrow CB$
$A \rightarrow AC$
$B \rightarrow 0$
$C \rightarrow 1$

$w = 10101101101$

Output: Can G derive w ?

$S \rightarrow AB \rightarrow CBB \rightarrow \dots \rightarrow w$

CYK, Earley's $\sim O(n^3)$

Valiant's Parser $O(n^\omega)$

Theorem: faster algorithms imply faster k-Clique!

[A-Backurs-VW Arxiv 15'] k-Clique based lower bounds

An important problem in Computational Biology

RNA Folding

Input: A sequence in $\{A, C, G, T\}^n$.



Matches:

A - T

C - G

Output: Maximum number of **non-crossing** matches

Best Algorithms are $\sim O(n^3)$

One reason to believe the new “k-Clique Conjecture”:

[Williams ICALP 04’]

Faster k-Clique implies faster Max-Cut

Best Exact Algs for Max-Cut: $2^{\omega n/3} < 1.73^n$

➔ What's next for "Hardness in P"?

- Better understanding of the conjectures
 - Classifying more problems
- More connections to Exact Algorithms

"More opportunities to say new things about fundamental problems"

CNF-SAT 1.9^n



Orthogonal Vectors $n^{1.9}$



“SETH-hard class”

...

LCS $n^{1.9}$

10-Dominating-Set $n^{9.9}$

Dynamic reachability $n^{0.9}$

Sparse Diameter $n^{1.9}$

Edit-Distance $n^{1.9}$

Frechet $n^{1.9}$

Max-Cut 1.7^n



k-clique $n^{0.7k}$



CFG Parsing $n^{2.3}$

RNA Folding $n^{2.3}$

your favorite NP-hard problem

Travelling Salesman 1.9^n

Set Cover 1.9^n

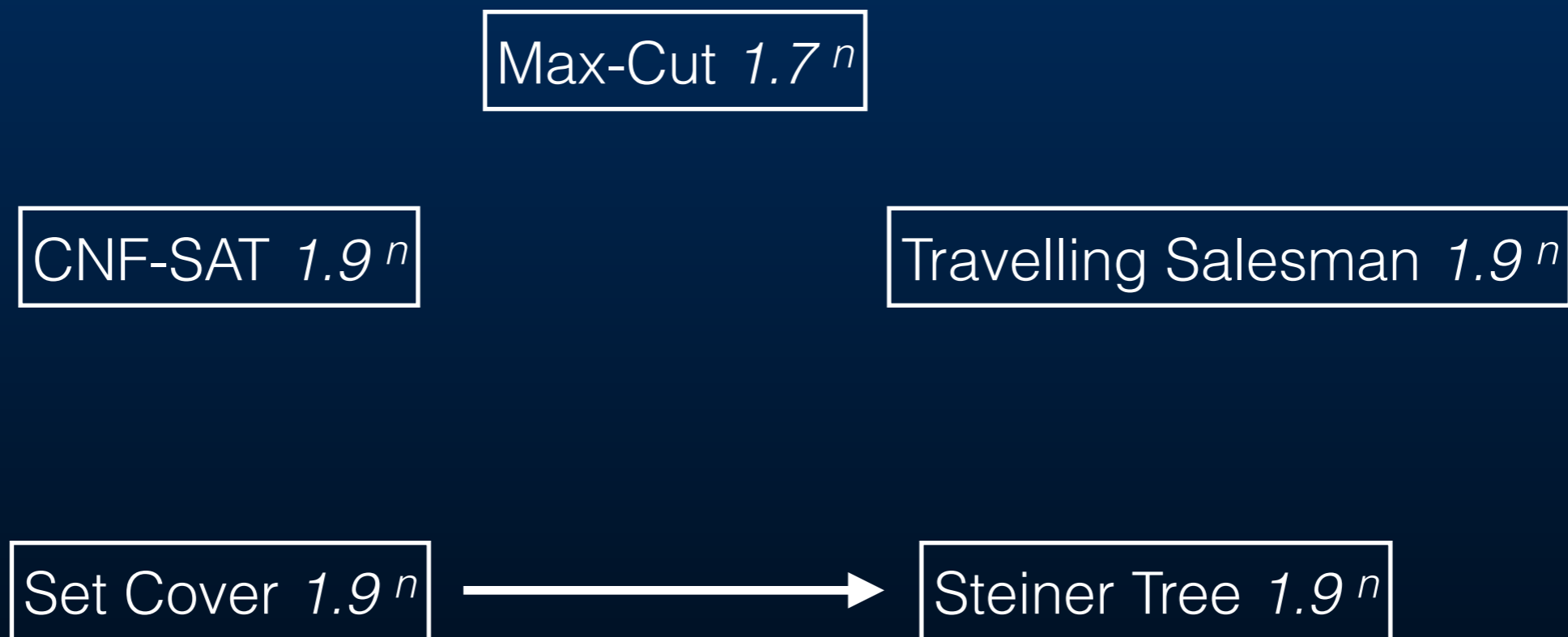


???



???

Tight Reductions within NP-hard problems?



[Cygan-Dell- Lokshantov-Marx-Nederlof-Okamoto-
Paturi-Saurabh-Wahlstrom CCC 12']

➔ What's next for "Hardness in P"?

- Better understanding of the conjectures
 - Classifying more problems
- More connections to Exact Algorithms
 - Fixed Parameter Tractability in P

"Apply the insights to P"

Parameterized Complexity:

Solve NP hard problems in $f(k) n^c$ time
on inputs of size n and some natural parameter k

[Giannopoulou - Mertzios - Niedermeier 15'] [A-VW-Wang 15']

*We should study Fixed Parameter Tractable algorithms
even for problems in P!*

Fixed Parameter Subquadratic

[A-VW-Wang 15']

Case study: Diameter in sparse graphs.
No subquadratic algorithm under SETH.

k = treewidth of G

Upper bound:

$$2^{O(k \log k)} n^{1+o(1)}$$

“Lower bound”:

$$2^{o(k)} n^{1.9}$$

refutes SETH

(the dependence on k is nearly tight!)

➔ What's next for "Hardness in P"?

- Better understanding of the conjectures
 - Classifying more problems
- More connections to Exact Algorithms
 - Fixed Parameter Tractability in P
 - Hardness of approximation

"Even more relevant lower bounds"

Best Approximation for Edit Distance in Subquadratic time?

Linear time approximations and heuristics are being used in practice

*Near-linear time **polylog** approximation is known*
[Andoni - Krauthgamer - Onak FOCS 10']

PCP Theorem in P?

➔ What's next for "Hardness in P"?

- Better understanding of the conjectures
 - Classifying more problems
- More connections to Exact Algorithms
 - Fixed Parameter Tractability in P
 - Hardness of approximation

Many more exciting questions:

- Barriers for shaving more log factors?
 - Average case hardness?
 - Quantum Algorithms?
 - Space Complexity?

Thanks for attending!