

CS156: The Calculus of Computation

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Chapter 3: First-Order Theories

First-Order Theories I

First-order theory T consists of

- ▶ Signature Σ_T - set of constant, function, and predicate symbols
- ▶ Set of axioms A_T - set of closed (no free variables) Σ_T -formulae

A Σ_T -formula is a formula constructed of constants, functions, and predicate symbols from Σ_T , and variables, logical connectives, and quantifiers.

The symbols of Σ_T are just symbols without prior meaning — the axioms of T provide their meaning.

First-Order Theories II

A Σ_T -formula F is valid in theory T (T -valid, also $T \models F$),
iff every interpretation I that satisfies the axioms of T ,

i.e. $I \models A$ for every $A \in A_T$ (T -interpretation)

also satisfies F ,

i.e. $I \models F$

A Σ_T -formula F is satisfiable in T (T -satisfiable), if there is a
 T -interpretation (i.e. satisfies all the axioms of T) that satisfies F

Two formulae F_1 and F_2 are equivalent in T (T -equivalent),

iff $T \models F_1 \leftrightarrow F_2$,

i.e. if for every T -interpretation I , $I \models F_1$ iff $I \models F_2$

Note:

- ▶ $I \models F$ stands for “ F true under interpretation I ”
- ▶ $T \models F$ stands for “ F is valid in theory T ”

Fragments of Theories

A fragment of theory T is a syntactically-restricted subset of formulae of the theory.

Example: a quantifier-free fragment of theory T is the set of quantifier-free formulae in T .

A theory T is decidable if $T \models F$ (T -validity) is decidable for every Σ_T -formula F ;

i.e., there is an algorithm that always terminate with “yes”, if F is T -valid, and “no”, if F is T -invalid.

A fragment of T is decidable if $T \models F$ is decidable for every Σ_T -formula F obeying the syntactic restriction.

Theory of Equality T_E I

Signature:

$$\Sigma_{=} : \{=, a, b, c, \dots, f, g, h, \dots, p, q, r, \dots\}$$

consists of

- ▶ $=$, a binary predicate, interpreted with meaning provided by axioms
- ▶ all constant, function, and predicate symbols

Axioms of T_E

1. $\forall x. x = x$ (reflexivity)
2. $\forall x, y. x = y \rightarrow y = x$ (symmetry)
3. $\forall x, y, z. x = y \wedge y = z \rightarrow x = z$ (transitivity)
4. for each positive integer n and n -ary function symbol f ,
 $\forall x_1, \dots, x_n, y_1, \dots, y_n. \bigwedge_i x_i = y_i$
 $\rightarrow f(x_1, \dots, x_n) = f(y_1, \dots, y_n)$ (function congruence)

Theory of Equality T_E II

5. for each positive integer n and n -ary predicate symbol p ,

$$\forall x_1, \dots, x_n, y_1, \dots, y_n. \bigwedge_i x_i = y_i \\ \rightarrow (p(x_1, \dots, x_n) \leftrightarrow p(y_1, \dots, y_n)) \text{ (predicate congruence)}$$

(function) and (predicate) are axiom schemata.

Example:

(function) for binary function f for $n = 2$:

$$\forall x_1, x_2, y_1, y_2. x_1 = y_1 \wedge x_2 = y_2 \rightarrow f(x_1, x_2) = f(y_1, y_2)$$

(predicate) for unary predicate p for $n = 1$:

$$\forall x, y. x = y \rightarrow (p(x) \leftrightarrow p(y))$$

Note: we omit “congruence” for brevity.

Decidability of T_E I

T_E is undecidable.

The quantifier-free fragment of T_E is decidable. Very efficient algorithm.

Semantic argument method can be used for T_E

Example: Prove

$$F : a = b \wedge b = c \rightarrow g(f(a), b) = g(f(c), a)$$

is T_E -valid.

Decidability of T_E II

Suppose not; then there exists a T_E -interpretation I such that $I \not\models F$. Then,

1. $I \not\models F$ assumption
2. $I \models a = b \wedge b = c$ 1, \rightarrow
3. $I \not\models g(f(a), b) = g(f(c), a)$ 1, \rightarrow
4. $I \models a = b$ 2, \wedge
5. $I \models b = c$ 2, \wedge
6. $I \models a = c$ 4, 5, (transitivity)
7. $I \models f(a) = f(c)$ 6, (function)
8. $I \models b = a$ 4, (symmetry)
9. $I \models g(f(a), b) = g(f(c), a)$ 7, 8, (function)
10. $I \models \perp$ 3, 9 contradictory

F is T_E -valid.

Natural Numbers and Integers

Natural numbers $\mathbb{N} = \{0, 1, 2, \dots\}$

Integers $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$

Three variations:

- ▶ Peano arithmetic T_{PA} : natural numbers with addition, multiplication, =
- ▶ Presburger arithmetic $T_{\mathbb{N}}$: natural numbers with addition, =
- ▶ Theory of integers $T_{\mathbb{Z}}$: integers with $+$, $-$, $>$, $=$, multiplication by constants

1. Peano Arithmetic T_{PA} (first-order arithmetic)

$$\Sigma_{PA} : \{0, 1, +, \cdot, =\}$$

Equality Axioms: (reflexivity), (symmetry), (transitivity)
two (function) for $+$, two (function) for \cdot .

And the axioms:

1. $\forall x. \neg(x + 1 = 0)$ (zero)
2. $\forall x, y. x + 1 = y + 1 \rightarrow x = y$ (successor)
3. $F[0] \wedge (\forall x. F[x] \rightarrow F[x + 1]) \rightarrow \forall x. F[x]$ (induction)
4. $\forall x. x + 0 = x$ (plus zero)
5. $\forall x, y. x + (y + 1) = (x + y) + 1$ (plus successor)
6. $\forall x. x \cdot 0 = 0$ (times zero)
7. $\forall x, y. x \cdot (y + 1) = x \cdot y + x$ (times successor)

Line 3 is an axiom schema.

Example: $3x + 5 = 2y$ can be written using Σ_{PA} as

$$x + x + x + 1 + 1 + 1 + 1 + 1 = y + y$$

Note: we have $>$ and \geq since

$$3x + 5 > 2y \quad \text{write as} \quad \exists z. z \neq 0 \wedge 3x + 5 = 2y + z$$

$$3x + 5 \geq 2y \quad \text{write as} \quad \exists z. 3x + 5 = 2y + z$$

Example:

Existence of pythagorean triples (F is T_{PA} -valid):

$$F : \exists x, y, z. x \neq 0 \wedge y \neq 0 \wedge z \neq 0 \wedge x \cdot x + y \cdot y = z \cdot z$$

Decidability of Peano Arithmetic

T_{PA} is undecidable. (Gödel, Turing, Post, Church)
The quantifier-free fragment of T_{PA} is undecidable.
(Matiyasevich, 1970)

Remark: Gödel's first incompleteness theorem

Peano arithmetic T_{PA} does not capture true arithmetic:

There exist closed Σ_{PA} -formulae representing valid propositions of number theory that are not T_{PA} -valid.

The reason: T_{PA} actually admits *nonstandard interpretations*.

For decidability: no multiplication

2. Presburger Arithmetic $T_{\mathbb{N}}$

Signature $\Sigma_{\mathbb{N}} : \{0, 1, +, =\}$

no multiplication!

Axioms of $T_{\mathbb{N}}$ (equality axioms, with 1-5):

1. $\forall x. \neg(x + 1 = 0)$ (zero)
2. $\forall x, y. x + 1 = y + 1 \rightarrow x = y$ (successor)
3. $F[0] \wedge (\forall x. F[x] \rightarrow F[x + 1]) \rightarrow \forall x. F[x]$ (induction)
4. $\forall x. x + 0 = x$ (plus zero)
5. $\forall x, y. x + (y + 1) = (x + y) + 1$ (plus successor)

Line 3 is an axiom schema.

$T_{\mathbb{N}}$ -satisfiability (and thus $T_{\mathbb{N}}$ -validity) is decidable
(Presburger, 1929)

3. Theory of Integers $T_{\mathbb{Z}}$

Signature:

$\Sigma_{\mathbb{Z}} : \{\dots, -2, -1, 0, 1, 2, \dots, -3\cdot, -2\cdot, 2\cdot, 3\cdot, \dots, +, -, >, =\}$

where

- ▶ $\dots, -2, -1, 0, 1, 2, \dots$ are constants
- ▶ $\dots, -3\cdot, -2\cdot, 2\cdot, 3\cdot, \dots$ are unary functions
(intended meaning: $2 \cdot x$ is $x + x$, $-3 \cdot x$ is $-x - x - x$)
- ▶ $+, -, >, =$ have the usual meanings.

Relation between $T_{\mathbb{Z}}$ and $T_{\mathbb{N}}$:

$T_{\mathbb{Z}}$ and $T_{\mathbb{N}}$ have the same expressiveness:

- ▶ For every $\Sigma_{\mathbb{Z}}$ -formula there is an equisatisfiable $\Sigma_{\mathbb{N}}$ -formula.
- ▶ For every $\Sigma_{\mathbb{N}}$ -formula there is an equisatisfiable $\Sigma_{\mathbb{Z}}$ -formula.

$\Sigma_{\mathbb{Z}}$ -formula F and $\Sigma_{\mathbb{N}}$ -formula G are *equisatisfiable* iff:

F is $T_{\mathbb{Z}}$ -satisfiable iff G is $T_{\mathbb{N}}$ -satisfiable

$\Sigma_{\mathbb{Z}}$ -formula to $\Sigma_{\mathbb{N}}$ -formula I

Example: consider the $\Sigma_{\mathbb{Z}}$ -formula

$$F_0 : \forall w, x. \exists y, z. x + 2y - z - 7 > -3w + 4.$$

Introduce two variables, v_p and v_n (range over the nonnegative integers) for each variable v (range over the integers) of F_0 :

$$F_1 : \forall w_p, w_n, x_p, x_n. \exists y_p, y_n, z_p, z_n. \\ (x_p - x_n) + 2(y_p - y_n) - (z_p - z_n) - 7 > -3(w_p - w_n) + 4$$

Eliminate $-$ by moving to the other side of $>$:

$$F_2 : \forall w_p, w_n, x_p, x_n. \exists y_p, y_n, z_p, z_n. \\ x_p + 2y_p + z_n + 3w_p > x_n + 2y_n + z_p + 7 + 3w_n + 4$$

$\Sigma_{\mathbb{Z}}$ -formula to $\Sigma_{\mathbb{N}}$ -formula II

Eliminate $>$ and numbers:

$$\forall w_p, w_n, x_p, x_n. \exists y_p, y_n, z_p, z_n. \exists u.$$

$$\begin{aligned} F_3 : \quad & \neg(u = 0) \wedge x_p + y_p + y_p + z_n + w_p + w_p + w_p \\ & = x_n + y_n + y_n + z_p + w_n + w_n + w_n + u \\ & + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 \end{aligned}$$

which is a $\Sigma_{\mathbb{N}}$ -formula equisatisfiable to F_0 .

$\Sigma_{\mathbb{Z}}$ -formula to $\Sigma_{\mathbb{N}}$ -formula III

Example: The $\Sigma_{\mathbb{N}}$ -formula

$$\forall x. \exists y. x = y + 1$$

is equisatisfiable to the $\Sigma_{\mathbb{Z}}$ -formula:

$$\forall x. x > -1 \rightarrow \exists y. y > -1 \wedge x = y + 1.$$

To decide $T_{\mathbb{Z}}$ -validity for a $\Sigma_{\mathbb{Z}}$ -formula F :

- ▶ transform $\neg F$ to an equisatisfiable $\Sigma_{\mathbb{N}}$ -formula $\neg G$,
- ▶ decide $T_{\mathbb{N}}$ -validity of G .

Rationals and Reals

Signatures:

$$\Sigma_{\mathbb{Q}} = \{0, 1, +, -, =, \geq\}$$

$$\Sigma_{\mathbb{R}} = \Sigma_{\mathbb{Q}} \cup \{\cdot\}$$

- ▶ Theory of Reals $T_{\mathbb{R}}$ (with multiplication)

$$x \cdot x = 2 \quad \Rightarrow \quad x = \pm\sqrt{2}$$

- ▶ Theory of Rationals $T_{\mathbb{Q}}$ (no multiplication)

$$\underbrace{2x}_{x+x} = 7 \quad \Rightarrow \quad x = \frac{7}{2}$$

Note: strict inequality okay; simply rewrite

$$x + y > z$$

as follows:

$$\neg(x + y = z) \wedge x + y \geq z$$

1. Theory of Reals $T_{\mathbb{R}}$

Signature:

$$\Sigma_{\mathbb{R}} : \{0, 1, +, -, \cdot, =, \geq\}$$

with multiplication. Axioms in text.

Example:

$$\forall a, b, c. b^2 - 4ac \geq 0 \leftrightarrow \exists x. ax^2 + bx + c = 0$$

is $T_{\mathbb{R}}$ -valid.

$T_{\mathbb{R}}$ is decidable (Tarski, 1930)
High time complexity

2. Theory of Rationals $T_{\mathbb{Q}}$

Signature:

$$\Sigma_{\mathbb{Q}} : \{0, 1, +, -, =, \geq\}$$

without multiplication. Axioms in text.

Rational coefficients are simple to express in $T_{\mathbb{Q}}$.

Example: Rewrite

$$\frac{1}{2}x + \frac{2}{3}y \geq 4$$

as the $\Sigma_{\mathbb{Q}}$ -formula

$$3x + 4y \geq 24$$

$T_{\mathbb{Q}}$ is decidable

Quantifier-free fragment of $T_{\mathbb{Q}}$ is efficiently decidable

Recursive Data Structures (RDS) I

Tuples of variables where the elements can be instances of the same structure: e.g., linked lists or trees.

1. Theory T_{CONS} (LISP-like lists)

Signature:

$$\Sigma_{\text{CONS}} : \{\text{cons}, \text{car}, \text{cdr}, \text{atom}, =\}$$

where

$\text{cons}(a, b)$ – list constructed by concatenating a and b

$\text{car}(x)$ – left projector of x : $\text{car}(\text{cons}(a, b)) = a$

$\text{cdr}(x)$ – right projector of x : $\text{cdr}(\text{cons}(a, b)) = b$

$\text{atom}(x)$ – true iff x is a single-element list

Note: an

atom is simply something that is not a cons. In this formulation, there is no NIL value.

Recursive Data Structures (RDS) II

Axioms:

1. The axioms of reflexivity, symmetry, and transitivity of =
2. Function Congruence axioms

$$\forall x_1, x_2, y_1, y_2. x_1 = x_2 \wedge y_1 = y_2 \rightarrow \text{cons}(x_1, y_1) = \text{cons}(x_2, y_2)$$

$$\forall x, y. x = y \rightarrow \text{car}(x) = \text{car}(y)$$

$$\forall x, y. x = y \rightarrow \text{cdr}(x) = \text{cdr}(y)$$

3. Predicate Congruence axiom

$$\forall x, y. x = y \rightarrow (\text{atom}(x) \leftrightarrow \text{atom}(y))$$

4. $\forall x, y. \text{car}(\text{cons}(x, y)) = x$ (left projection)
5. $\forall x, y. \text{cdr}(\text{cons}(x, y)) = y$ (right projection)
6. $\forall x. \neg \text{atom}(x) \rightarrow \text{cons}(\text{car}(x), \text{cdr}(x)) = x$ (construction)
7. $\forall x, y. \neg \text{atom}(\text{cons}(x, y))$ (atom)

Note: the behavior of car and cons on atoms is not specified.

T_{cons} is undecidable

Quantifier-free fragment of T_{cons} is efficiently decidable

Lists with equality

2. Theory T_{cons}^E (lists with equality)

$$T_{\text{cons}}^E = T_E \cup T_{\text{cons}}$$

Signature:

$$\Sigma_E \cup \Sigma_{\text{cons}}$$

(this includes uninterpreted constants, functions, and predicates)

Axioms: union of the axioms of T_E and T_{cons}

T_{cons}^E is undecidable

Quantifier-free fragment of T_{cons}^E is efficiently decidable

Example: The Σ_{cons}^E -formula

$$F : \begin{aligned} & \text{car}(x) = \text{car}(y) \wedge \text{cdr}(x) = \text{cdr}(y) \wedge \neg \text{atom}(x) \wedge \neg \text{atom}(y) \\ & \rightarrow f(x) = f(y) \end{aligned}$$

is T_{cons}^E -valid.

Suppose not; then there exists a T_{cons}^E -interpretation I such that $I \not\models F$. Then,

- | | | |
|-----|---|--------------------------|
| 1. | $I \not\models F$ | assumption |
| 2. | $I \models \text{car}(x) = \text{car}(y)$ | 1, \rightarrow, \wedge |
| 3. | $I \models \text{cdr}(x) = \text{cdr}(y)$ | 1, \rightarrow, \wedge |
| 4. | $I \models \neg \text{atom}(x)$ | 1, \rightarrow, \wedge |
| 5. | $I \models \neg \text{atom}(y)$ | 1, \rightarrow, \wedge |
| 6. | $I \not\models f(x) = f(y)$ | 1, \rightarrow |
| 7. | $I \models \text{cons}(\text{car}(x), \text{cdr}(x)) = \text{cons}(\text{car}(y), \text{cdr}(y))$ | 2, 3, (function) |
| 8. | $I \models \text{cons}(\text{car}(x), \text{cdr}(x)) = a$ | 4, (construction) |
| 9. | $I \models \text{cons}(\text{car}(y), \text{cdr}(y)) = b$ | 5, (construction) |
| 10. | $I \models a = b$ | 7, 8, 9, (transitivity) |
| 11. | $I \models f(x) = f(y)$ | 10, (function) |

Lines 6 and 11 are contradictory, so our assumption that $I \not\models F$ must be wrong. Therefore, F is T_{cons}^E -valid.

Theory of Arrays T_A

Signature:

$$\Sigma_A : \{ \cdot[\cdot], \cdot\langle \cdot \triangleleft \cdot \rangle, = \}$$

where

- ▶ $a[i]$ binary function –
read array a at index i (“read(a, i)”)
- ▶ $a\langle i \triangleleft v \rangle$ ternary function –
write value v to index i of array a (“write(a, i, v)”)

Axioms

1. the axioms of (reflexivity), (symmetry), and (transitivity) of T_E
2. $\forall a, i, j. i = j \rightarrow a[i] = a[j]$ (array congruence)
3. $\forall a, v, i, j. i = j \rightarrow a\langle i \triangleleft v \rangle[j] = v$ (read-over-write 1)
4. $\forall a, v, i, j. i \neq j \rightarrow a\langle i \triangleleft v \rangle[j] = a[j]$ (read-over-write 2)

Note: = is only defined for array elements

$$F : a[i] = e \rightarrow a\langle i \triangleleft e \rangle = a$$

not T_A -valid, but

$$F' : a[i] = e \rightarrow \forall j. a\langle i \triangleleft e \rangle[j] = a[j] ,$$

is T_A -valid.

Also

$$a = b \rightarrow a[i] = b[i]$$

is not T_A -valid: We have only axiomatized a restricted congruence.

T_A is undecidable

Quantifier-free fragment of T_A is decidable

2. Theory of Arrays T_A^- (with extensionality)

Signature and axioms of T_A^- are the same as T_A , with one additional axiom

$$\forall a, b. (\forall i. a[i] = b[i]) \leftrightarrow a = b \quad (\text{extensionality})$$

Example:

$$F : a[i] = e \rightarrow a\langle i \triangleleft e \rangle = a$$

is T_A^- -valid.

T_A^- is undecidable
Quantifier-free fragment of T_A^- is decidable

First-Order Theories

	Theory	Decidable	QFF Dec.
T_E	Equality	—	✓
T_{PA}	Peano Arithmetic	—	—
$T_{\mathbb{N}}$	Presburger Arithmetic	✓	✓
$T_{\mathbb{Z}}$	Linear Integer Arithmetic	✓	✓
$T_{\mathbb{R}}$	Real Arithmetic	✓	✓
$T_{\mathbb{Q}}$	Linear Rationals	✓	✓
T_{cons}	Lists	—	✓
T_{cons}^E	Lists with Equality	—	✓

Combination of Theories

How do we show that

$$1 \leq x \wedge x \leq 2 \wedge f(x) \neq f(1) \wedge f(x) \neq f(2)$$

is $(T_E \cup T_{\mathbb{Z}})$ -unsatisfiable?

Or how do we prove properties about
an array of integers, or
a list of reals ...?

Given theories T_1 and T_2 such that

$$\Sigma_1 \cap \Sigma_2 = \{=\}$$

The combined theory $T_1 \cup T_2$ has

- ▶ signature $\Sigma_1 \cup \Sigma_2$
- ▶ axioms $A_1 \cup A_2$

Nelson & Oppen showed that,
if

- ▶ satisfiability of the quantifier-free fragment (qff) of T_1 is decidable,
 - ▶ satisfiability of qff of T_2 is decidable, and
 - ▶ certain technical simple requirements are met,
- then satisfiability of qff of $T_1 \cup T_2$ is decidable.